

Into the Unknown: Assigning Reviewers to Papers with Uncertain Affinities



Cyrus Cousins



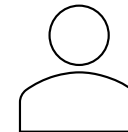
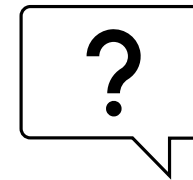
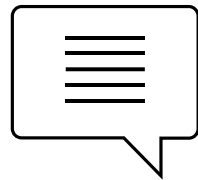
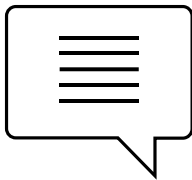
Justin Payan



Yair Zick




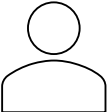



Peer Review

Papers must be reviewed by suitable reviewers


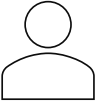

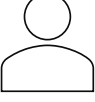

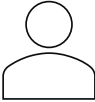


Reviewer Assignment Problem







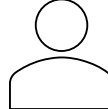

Given *affinity scores*

				
	5	4	8	4
	9	6	1	8
	3	4	6	7

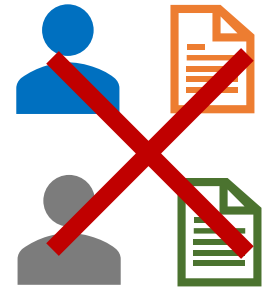
Paper
Requirements:

	2	
	2	
	3	

Reviewer Load
Bounds:



	≤ 2	
	≤ 3	
	≤ 1	
	≤ 2	

Conflicts
of
interest:



Goal:
Allocation of reviewers to papers with high affinity

What are affinity scores?

For given *reviewer-paper pair* →   5

the affinity score measures:

Reviewer
interest



Reviewer
expertise





As proxies for:

Predicted review quality



What are affinity scores?

For given *reviewer-paper pair* →   5

Reviewer
interest

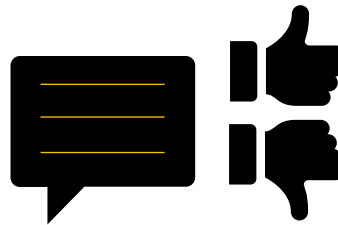


Reviewer
expertise



As proxies for:




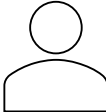



Predicted review quality

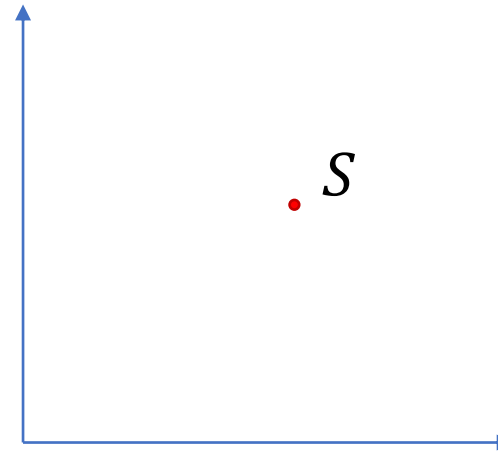


Hard to
measure/predict

Reviewer Assignment Under Uncertainty (RAU)




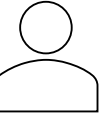



Compute an *uncertainty set* \mathcal{S}
s.t. true, unknown affinity scores S are contained in \mathcal{S} with prob. $1 - \delta$

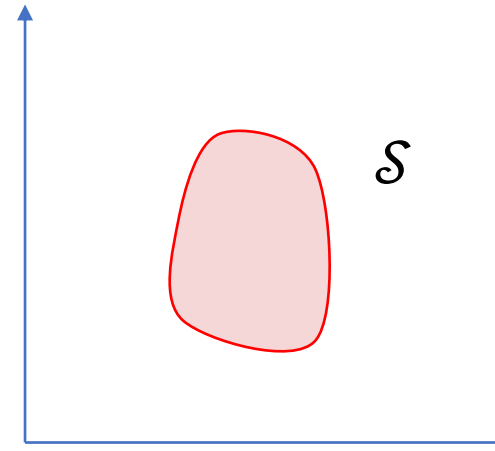
				
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Reviewer Assignment Under Uncertainty (RAU)




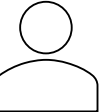



Compute an *uncertainty set* \mathcal{S}
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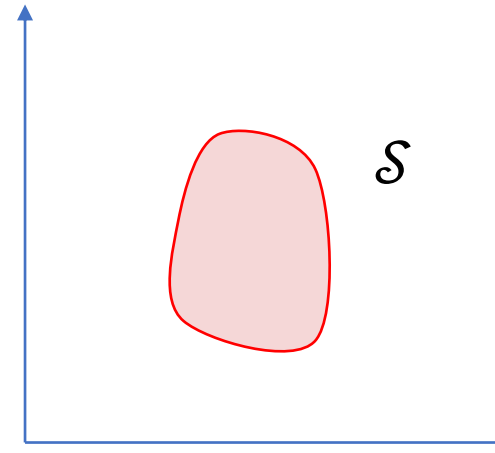
				
	$\cong 5$	$\cong 4$	$\cong 8$	$\cong 4$
	$\cong 9$	$\cong 6$	$\cong 1$	$\cong 8$
	$\cong 3$	$\cong 4$	$\cong 6$	$\cong 7$



Reviewer Assignment Under Uncertainty (RAU)

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	$\cong 5$	$\cong 4$	$\cong 8$	$\cong 4$
	$\cong 9$	$\cong 6$	$\cong 1$	$\cong 8$
	$\cong 3$	$\cong 4$	$\cong 6$	$\cong 7$



Maximize the worst-case welfare over the *uncertainty set*

Reviewer Assignment Under Uncertainty (RAU)

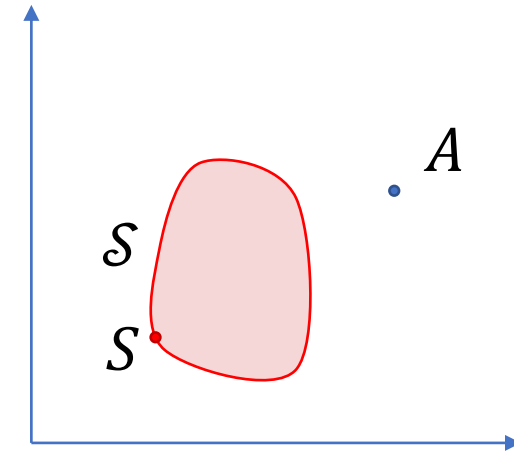
Papers P and reviewers R

Valid assignments $\mathcal{A} \subseteq \{0, 1\}^{|P| \times |R|}$

Affinity score uncertainty set $\mathcal{S} \subseteq \mathbb{R}_+^{|P| \times |R|}$

$$USW(A, S) = \frac{1}{|P|} \sum_{p \in P} \sum_{r \in R} A_{pr} S_{pr}$$

Objective: $\max_{A \in \mathcal{A}} \min_{S \in \mathcal{S}} USW(A, S)$



Reviewer Assignment Under Uncertainty (RAU)

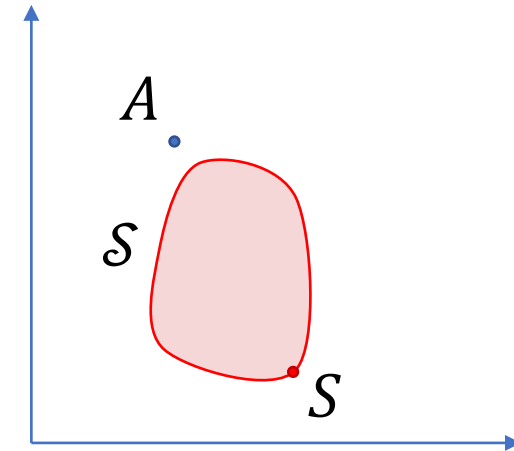
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Why Maximin?

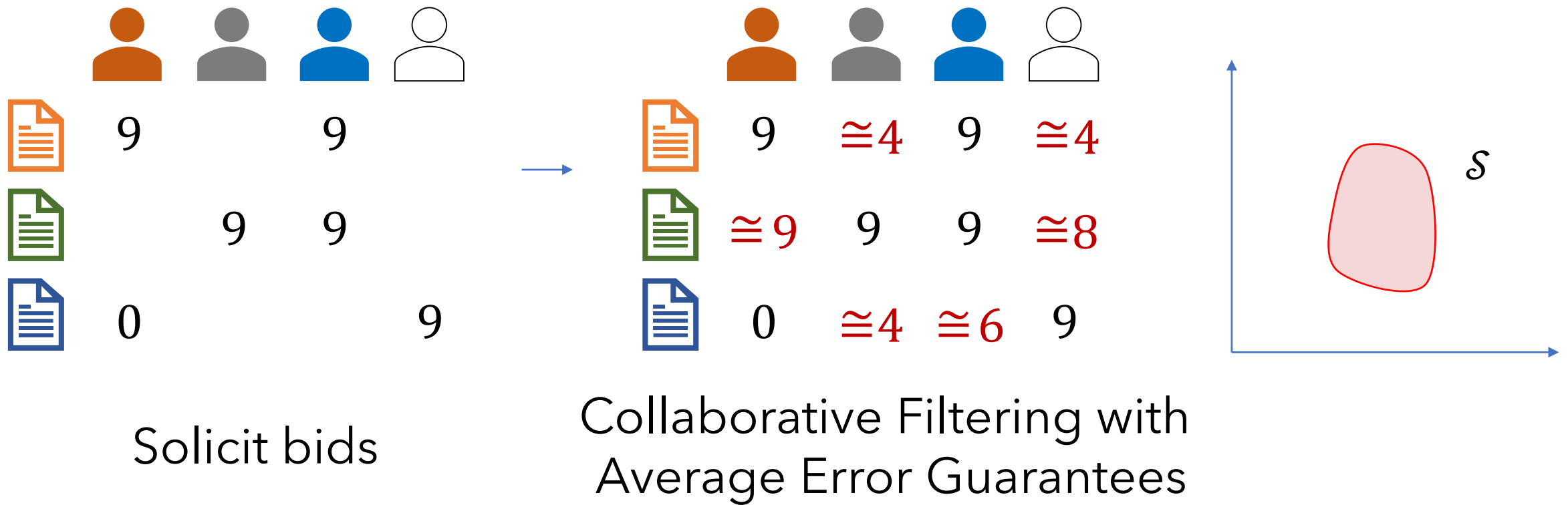
Why Maximin?

- 1 May not always know expected affinities
- 2 Theory and experiments show maximin corresponds to *true* welfare

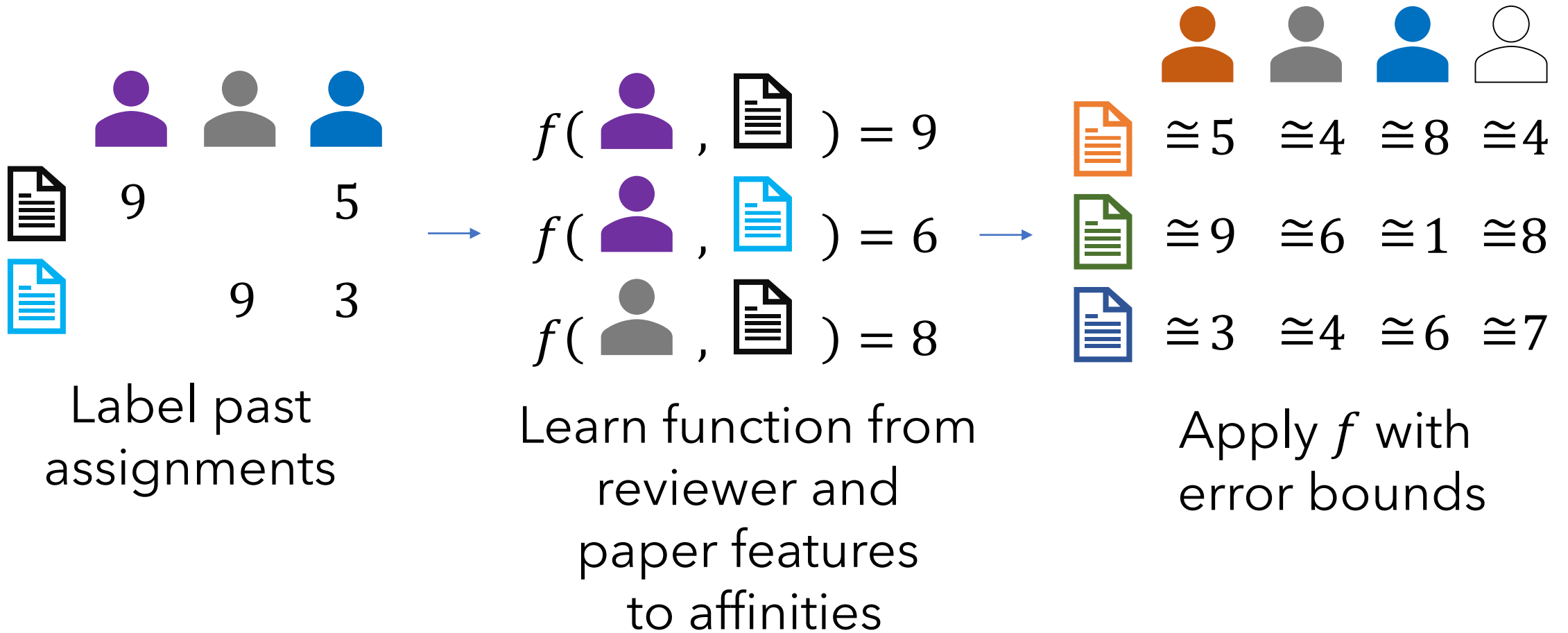


Implementing RAU

Implementing RAU: Bid Prediction

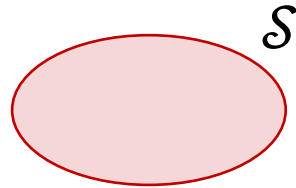


Implementing RAU: Review Quality Prediction





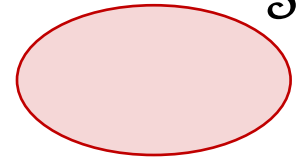
Transductive Predictors
(CF on Bids)



\mathcal{S}

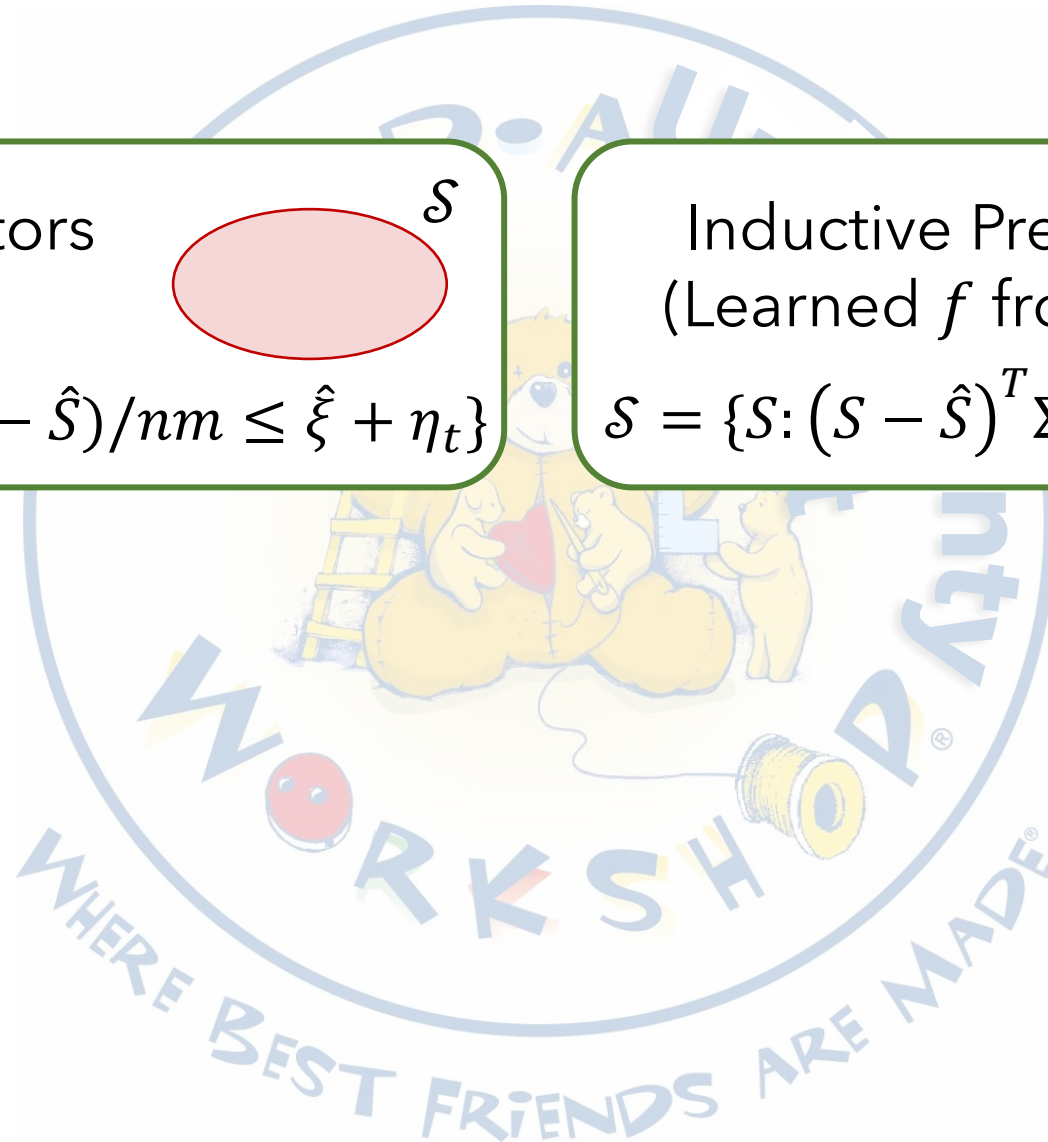
$$\mathcal{S} = \{S: (S - \hat{S})^T \Sigma^{-1} (S - \hat{S}) / nm \leq \hat{\xi} + \eta_t\}$$

Inductive Predictors
(Learned f from Data)

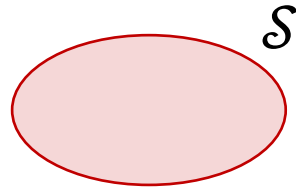


\mathcal{S}

$$\mathcal{S} = \{S: (S - \hat{S})^T \Sigma^{-1} (S - \hat{S}) / nm \leq \hat{\xi} + \eta_i\}$$



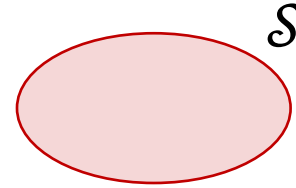
Transductive Predictors
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\mathcal{S}

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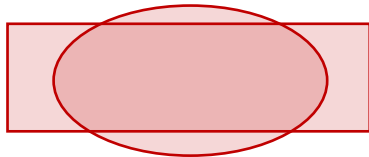
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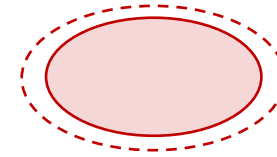
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Intersecting Uncertainty Sets



$$p(S \in \mathcal{S}) = 1 - (\delta_1 + \delta_2)$$

Expanding Uncertainty Sets



A man with a mustache is blowing a bubble with gum. He is wearing a dark blue vest over a white shirt. He is standing in front of a whiteboard covered in many colorful sticky notes (yellow, blue, pink). The text "Solving RAU" is overlaid on the image.

Solving RAU

RAU is NP-hard

Theorem 1:
RAU is NP-hard for convex uncertainty sets \mathcal{S}

RAU is NP-hard

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RAU is NP-hard for convex uncertainty sets \mathcal{S}

Nonlinear objective


$$\max_{A \in \mathcal{A}} \min_{S \in \mathcal{S}} USW(A, S)$$

Robust Reviewer Assignment (RRA)

Relax discrete allocations \rightarrow continuous

$$\mathcal{A} \subseteq \{0, 1\}^{m \times n} \rightarrow \tilde{\mathcal{A}} \subseteq [0, 1]^{m \times n}$$

Projected subgradient-ascent optimization

Solve $\max_{A \in \tilde{\mathcal{A}}} \min_{S \in \mathcal{S}} USW(A, S)$
by stepping in $\partial_A \min_{S \in \mathcal{S}} USW(A, S)$
and projecting back to $\tilde{\mathcal{A}}$

Randomized rounding for discrete solution

Round $A \in \tilde{\mathcal{A}}$ to $A' \in \mathcal{A}$

RRA Guarantees

Define

\tilde{A} : the continuous maximin solution

A' : the rounding of \tilde{A}

S^* : the true affinity score matrix

A^* : the true optimal allocation

L : \mathcal{L}_1 -diameter of \mathcal{S}

δ : $\mathbb{P}(S^* \in \mathcal{S}) \geq 1 - \delta$

RRA Guarantees

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Integrality Gap

$$\begin{aligned}\mathbb{E}_{A'} \|A' - \tilde{A}\|_1 &= 2 \left(\|A'\|_1 - \|\tilde{A}\|_1 \right) \\ &= |P||R| - 2 \left\| \frac{1}{2} - \tilde{A} \right\|_1\end{aligned}$$

RRA Guarantees

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True Welfare Gap

$$\mathbb{P} \left(USW(A^*, S^*) - \mathbb{E}_{A'} USW(A', S^*) > \frac{L}{|P|} \right) < \delta$$

RRA Guarantees

Integrality Gap

We may have to round heavily...

But the maximin objective still ensures high *true* welfare!

$< \delta$

ICLR Experiments

			100 * Adversarial USW		100 * Average USW	
Year	# Revs	# Papers	ILP	RRA	ILP	RRA
2018	1657	546	17	16	179	160
2019	2620	851	22	27	184	161
2020	4123	1327	17	23	187	166
2021	4662	1557	23	33	192	174
2022	5023	1576	28	38	191	172
Average			21.4	27.4 (+28%)	187	167 (-11%)

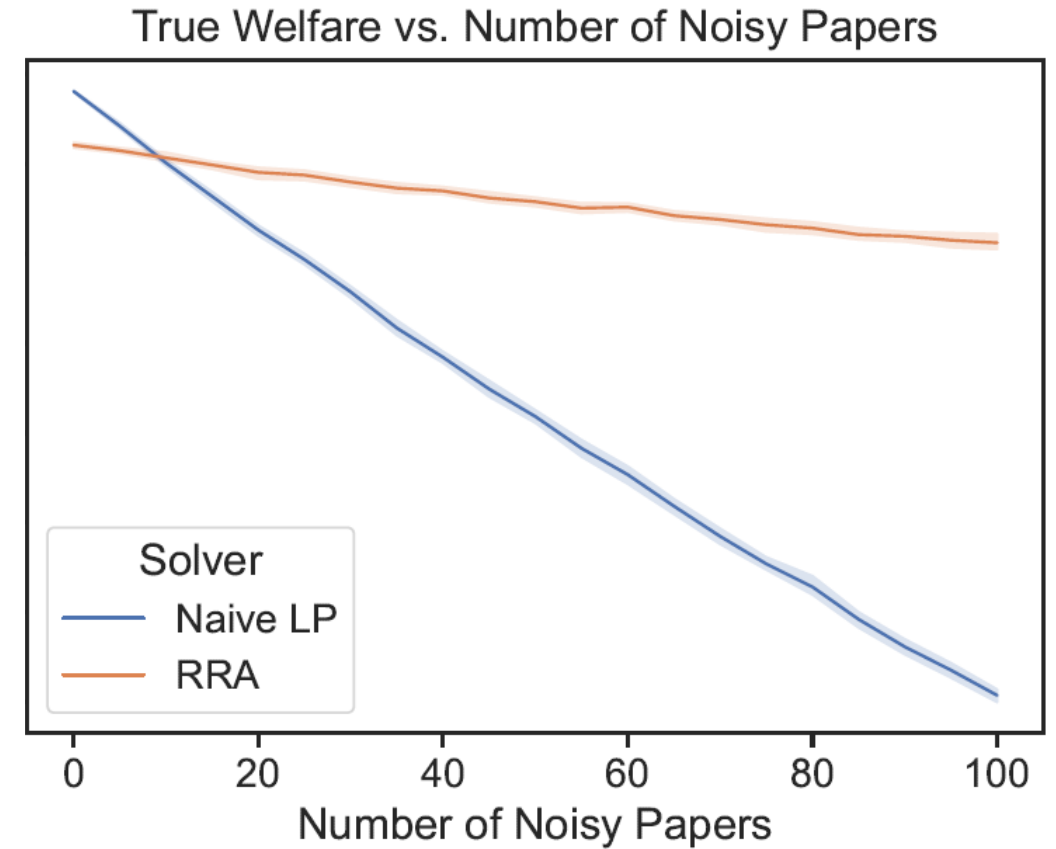
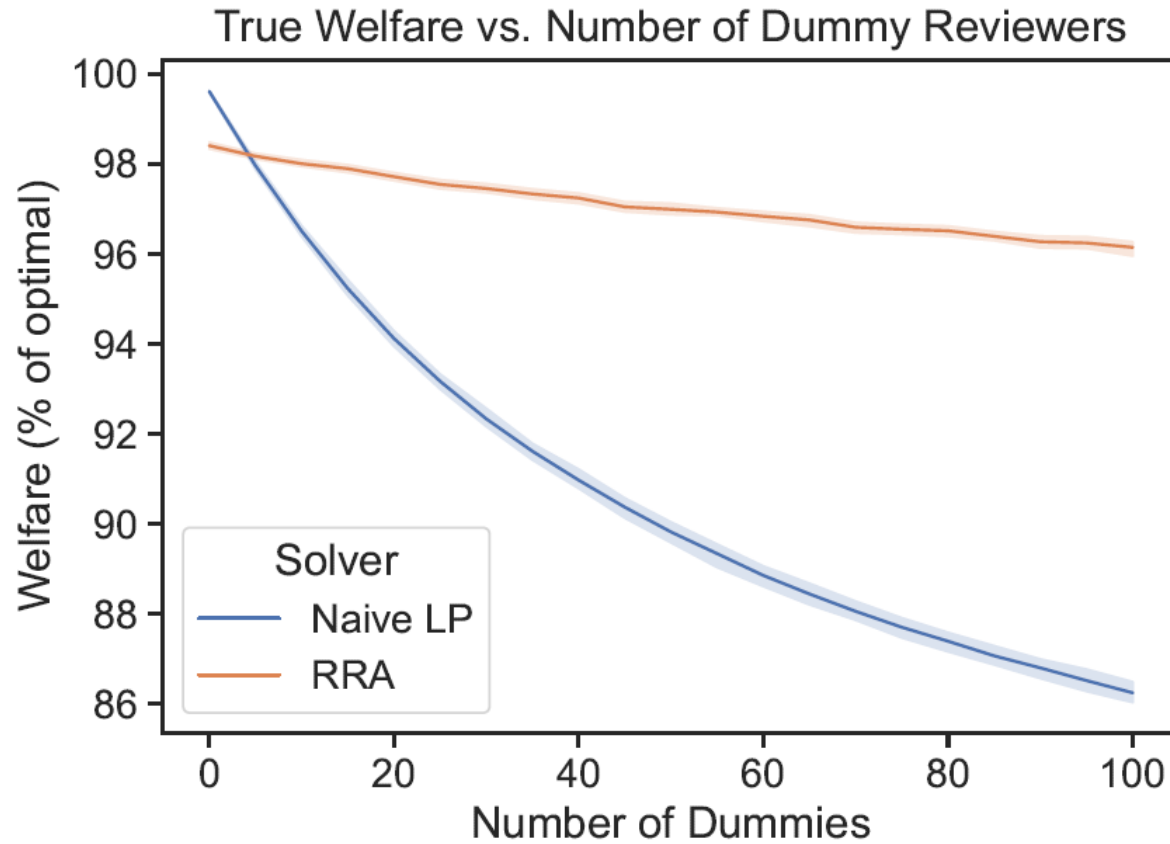
USW numbers are
divided by # Papers

ICLR Experiments

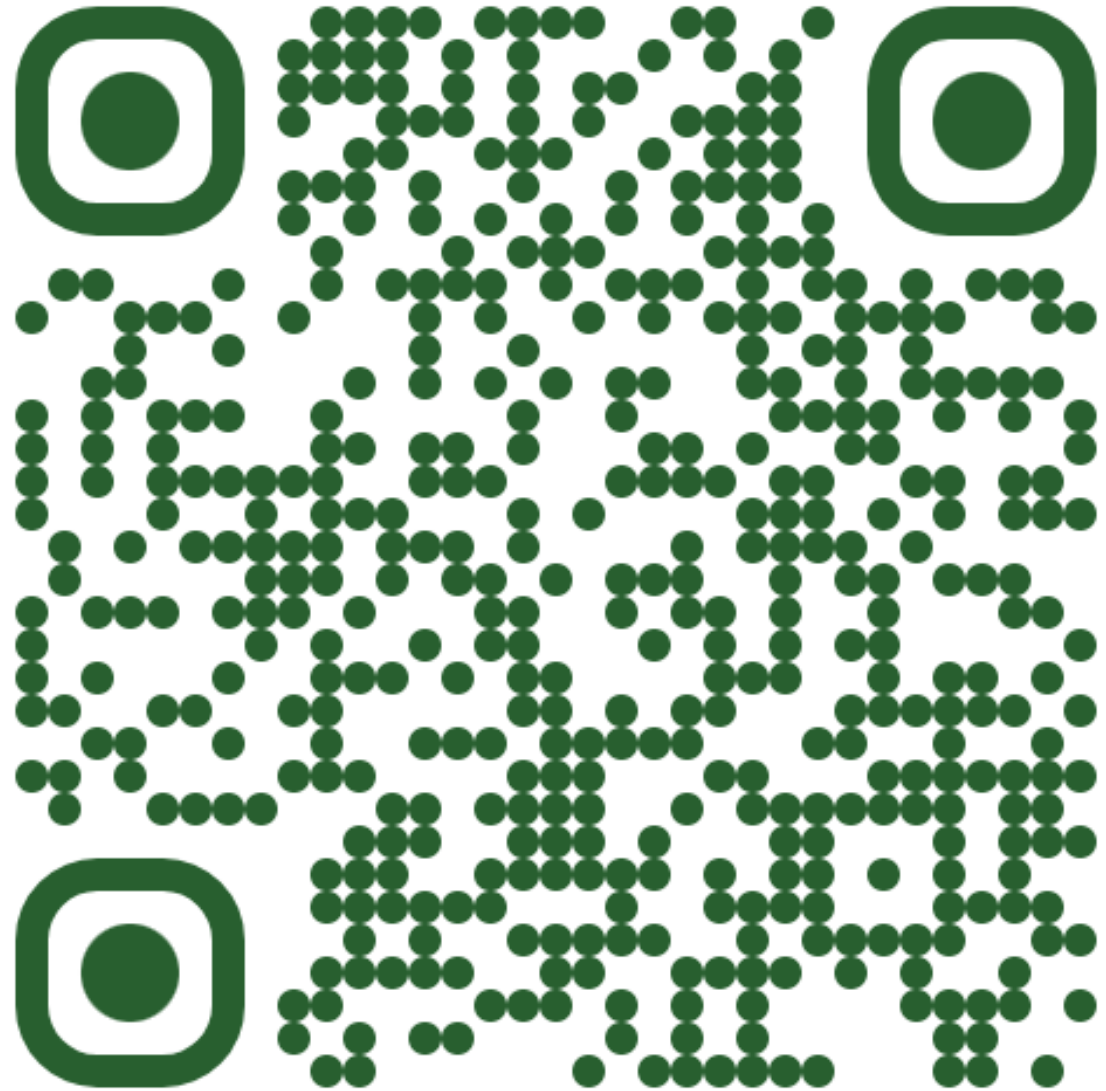
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USW numbers are
divided by # Papers

True Welfare under Noise



https://t.ly/g_sLd



Thanks!

Email:
jpayan@cs.umass.edu

Robust Reviewer Assignment (RRA)

What about general convex uncertainty sets \mathcal{S} ?

Relax, then apply projected subgradient ascent against an “adversary”

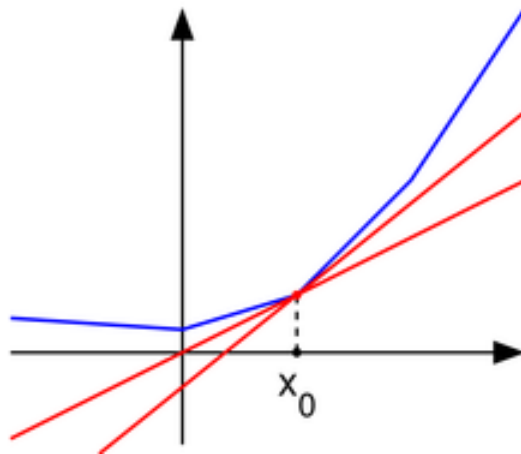
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Gradient ascent for
non-differentiable functions



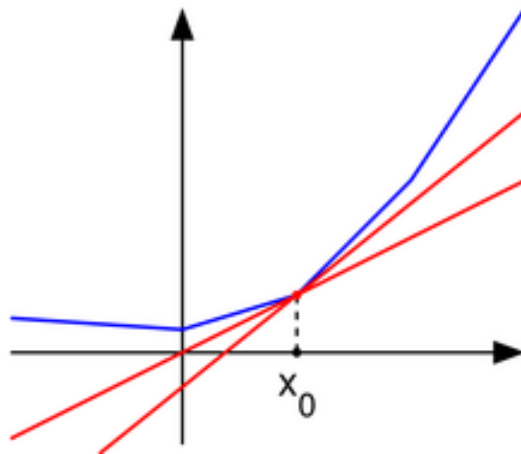
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At each step, ensure
constraints are satisfied

Gradient ascent for
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Robust Reviewer Assignment (RRA)

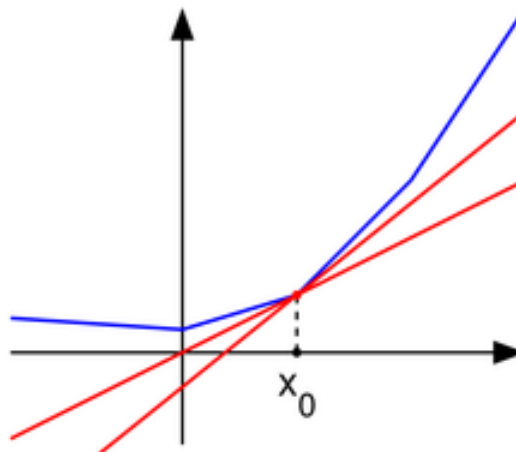
What about general convex uncertainty sets \mathcal{S} ?

Relax, then apply projected subgradient ascent against an “adversary”

At each step, ensure constraints are satisfied

Gradient ascent for non-differentiable functions

It's a maximin problem



Robust Reviewer Assignment (RRA)

What about general convex uncertainty sets \mathcal{S} ?

Relax, then apply projected subgradient ascent against an “adversary”

Converges in poly-time if projection and adversary are efficient

Robust Reviewer Assignment (RRA)

Randomized rounding for discrete solution

Round $A \in \tilde{\mathcal{A}}$ to $A' \in \mathcal{A}$

Robust Reviewer Assignment (RRA)

Randomized rounding for discrete solution

Round $A \in \tilde{\mathcal{A}}$ to $A' \in \mathcal{A}$

Generalization of Birkhoff von Neumann decomposition

Maintains constraints of $\tilde{\mathcal{A}}$ and \mathcal{A}

$$\mathbb{E}[A'] = A$$

$$\min_{S \in \mathcal{S}} USW(A', S) \leq \min_{S \in \mathcal{S}} USW(A, S)$$

ICLR Experiments

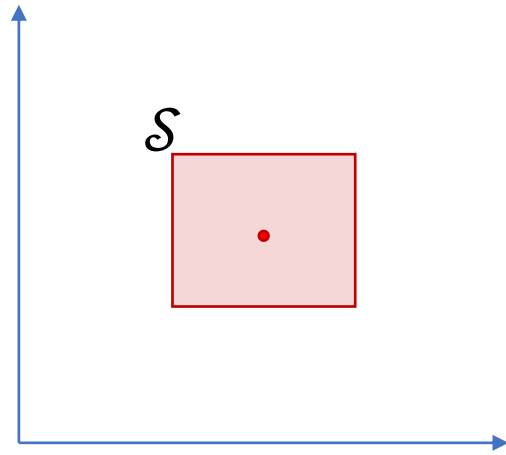
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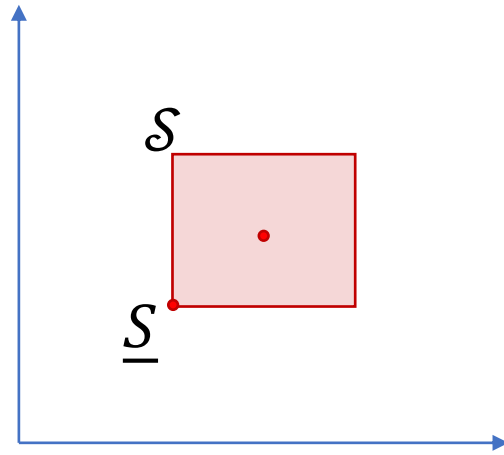
Simple Uncertainty Sets

Independent
Bounding Boxes



Simple Uncertainty Sets

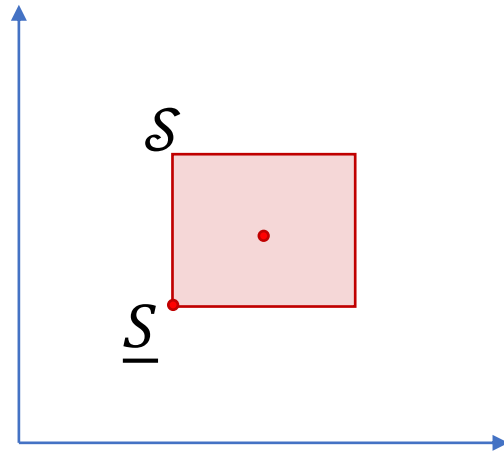
Independent
Bounding Boxes



Optimize
assuming \underline{S}

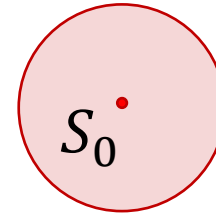
Simple Uncertainty Sets

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Sphere

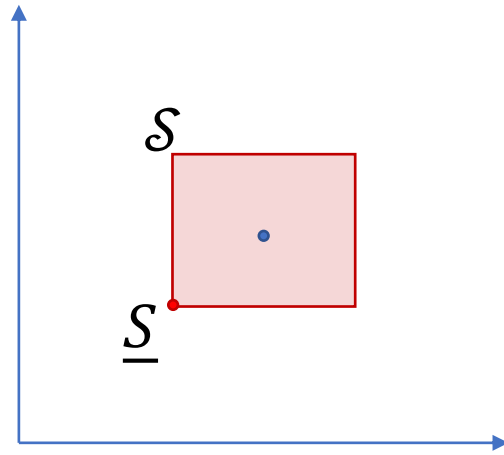
$$\mathcal{S} = \mathcal{B}_\epsilon(S_0)$$



Optimize
assuming \underline{S}

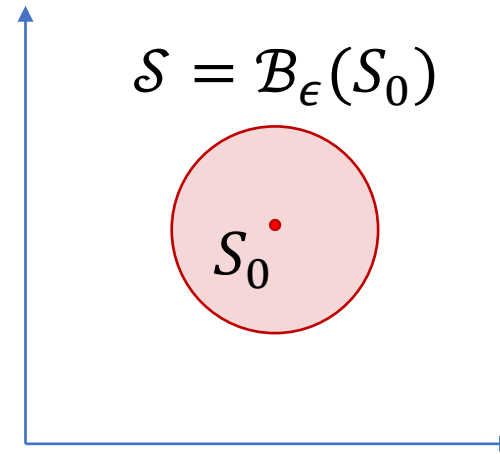
Simple Uncertainty Sets

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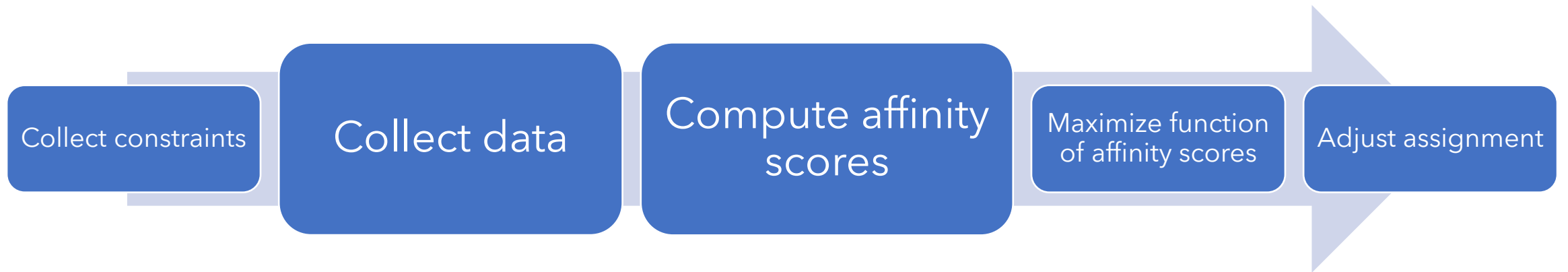


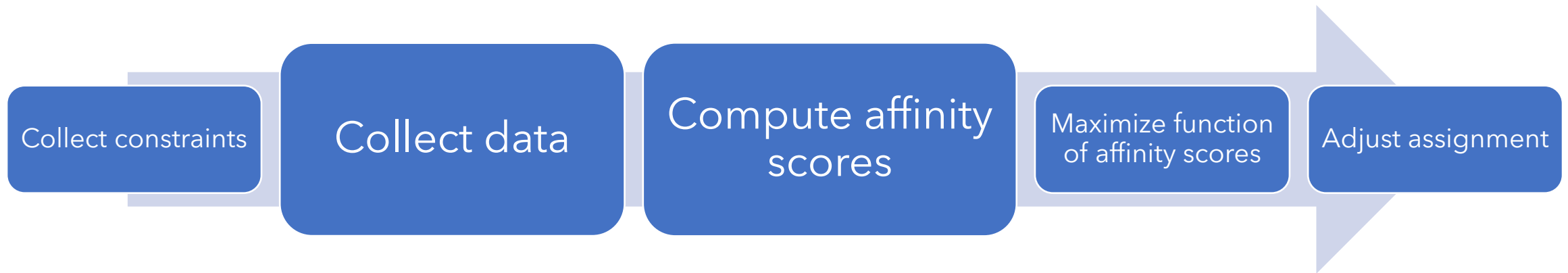
Optimize
assuming \underline{S}

Sphere



Optimize
assuming S_0



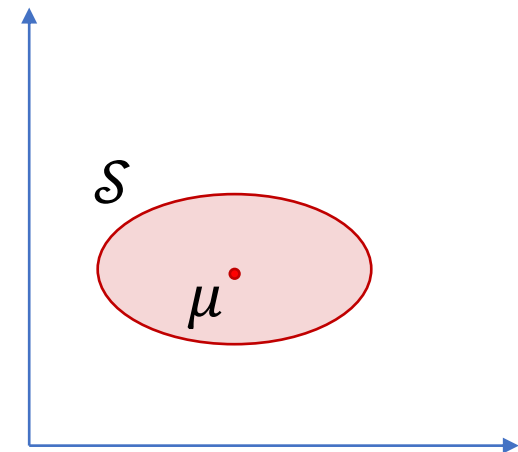


$$\text{Base Aggregated Score} = \begin{cases} \frac{1}{4}\text{TPMS} + \frac{1}{4}\text{ACL} + \frac{1}{2}\text{SAM} & \text{if all data is present} \\ \frac{1}{2}\text{ACL} + \frac{1}{2}\text{SAM} & \text{if TPMS is missing} \\ \frac{1}{2}\text{TPMS} + \frac{1}{2}\text{SAM} & \text{if ACL is missing} \\ \text{SAM} & \text{if both ACL and TPMS are missing} \end{cases}$$

Compute \mathcal{S} as a $\tau\%$ CI of $\mathcal{N}(\mu, \Sigma)$

μ is the base scores

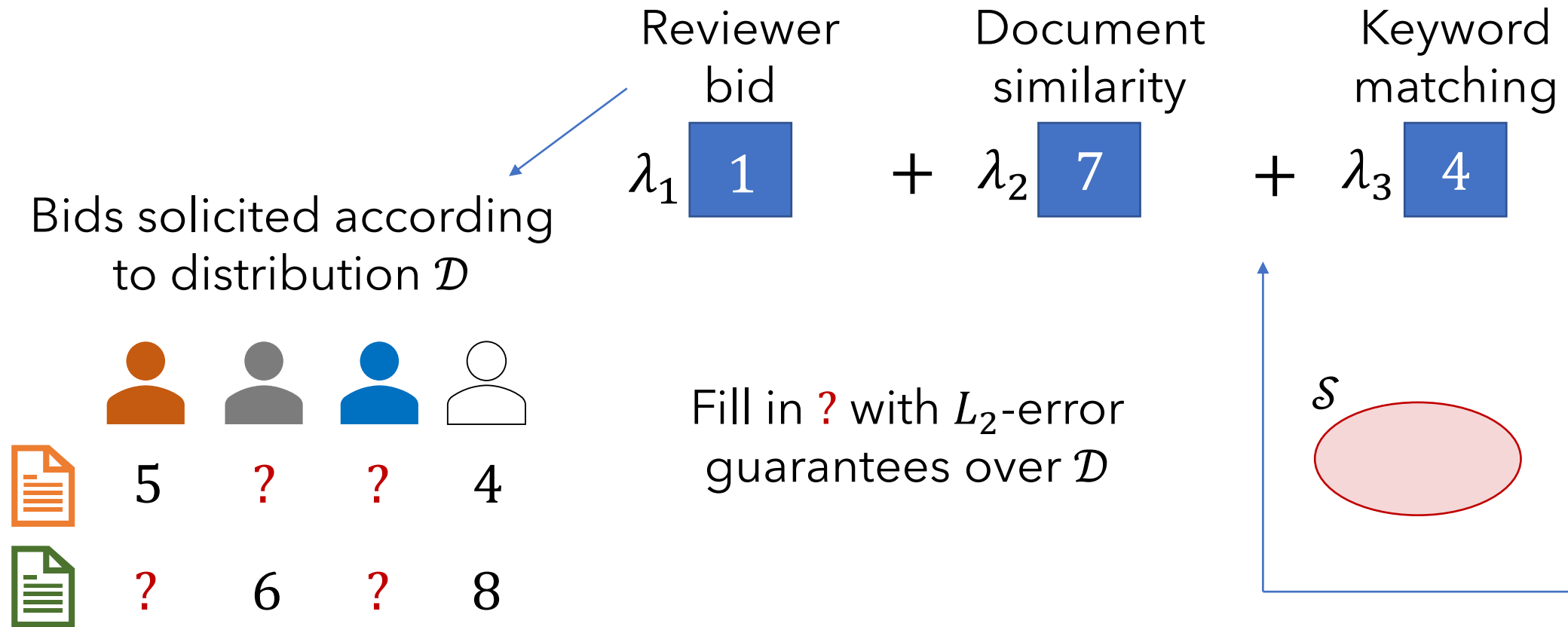
Σ is diagonal, based on number of missing sources





Extensions (and Related Problems)

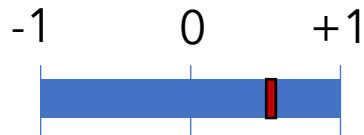
Learning Bids



Assign for Predicted Quality

Predict important review aspects

Agreement with final decision (d)



Probability of covering aspects ($p(a_i)$)



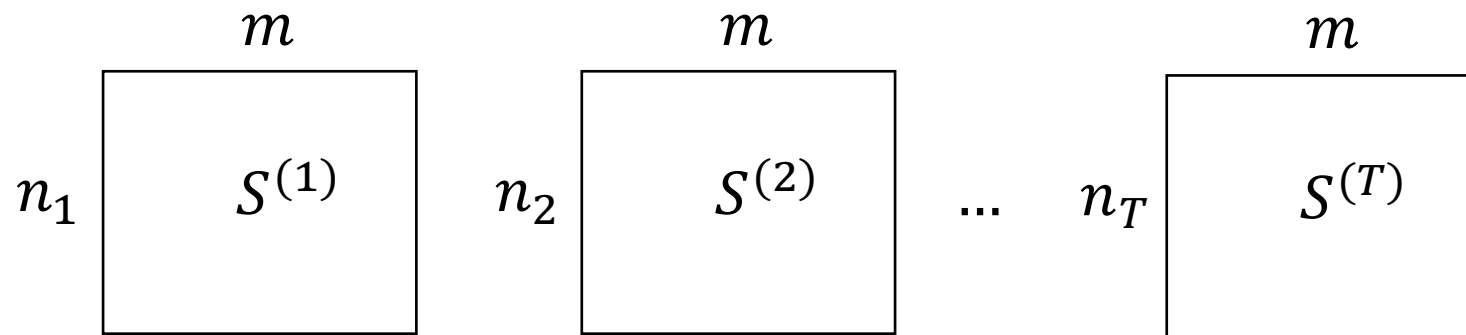
Value of assigning reviewer set \mathcal{C} to paper p is

$$\lambda_1 \sum_{r \in \mathcal{C}} d_{rp} + \lambda_2 \sum_i \max_{r \in \mathcal{C}} p(a_i | r, p)$$

Constrained allocation of indivisible resources with submodular valuations

Online Reviewer Assignment

m reviewers available, must review n_t papers in month t



Goal: $\max_{A^{(1)}, A^{(2)}, \dots, A^{(T)}} \sum_t USW(A^{(t)}, S^{(t)})$ subj. to

$$\sum_t \sum_p A_{pr}^{(t)} \leq U_r \quad \text{for all } r$$

$$\sum_p A_{pr}^{(t)} \leq U_r^{(t)} \quad \text{for all } r, t$$

Improving RRA

Precise trade-offs in randomized rounding with constraints

Randomized rounding for discrete solution

Round $A \in \tilde{\mathcal{A}}$ to $A' \in \mathcal{A}$

Generalization of Birkhoff von Neumann decomposition

Maintains constraints of $\tilde{\mathcal{A}}$ and \mathcal{A}

$$\mathbb{E}[A'] = A$$

Can replace
this with
guarantees
on USW?

Uncertainty: Nowhere to be Found

Cannot trust high affinity scores

Low scores are too pessimistic
(esp. with missing data)



Outline



Step through a typical conference workflow



Introduce RRA, a framework that accounts for uncertainty




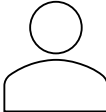



Also allows us to use new affinity score estimation methods

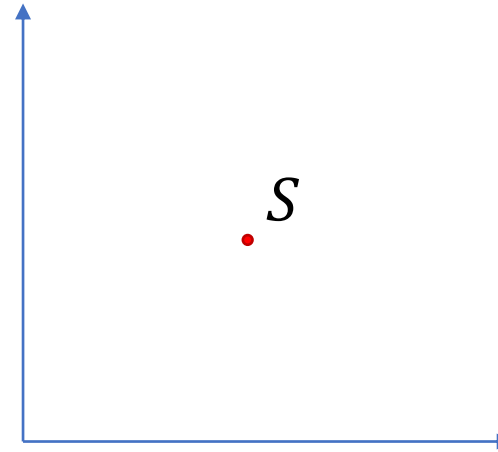
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Compute an *uncertainty set* \mathcal{S} containing true, unknown affinity scores S

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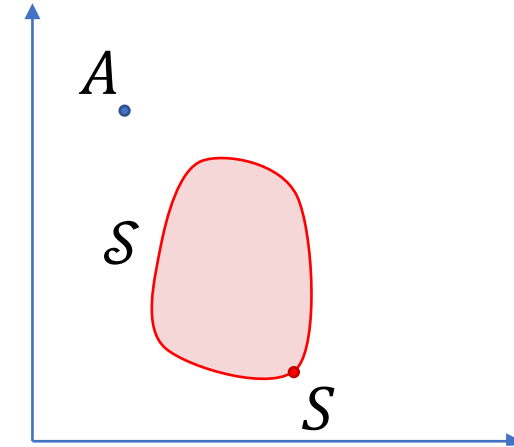
What is RAU, really?

What is RAU, really?

Maximin USW is a lower bound for *true* USW

Flexibly trades off between
certainty and high expected value

Approaches original problem as $\mathcal{S} \rightarrow S$



$$\max_{A \in \mathcal{A}} \min_{S \in \mathcal{S}} USW(A, S)$$

Outline



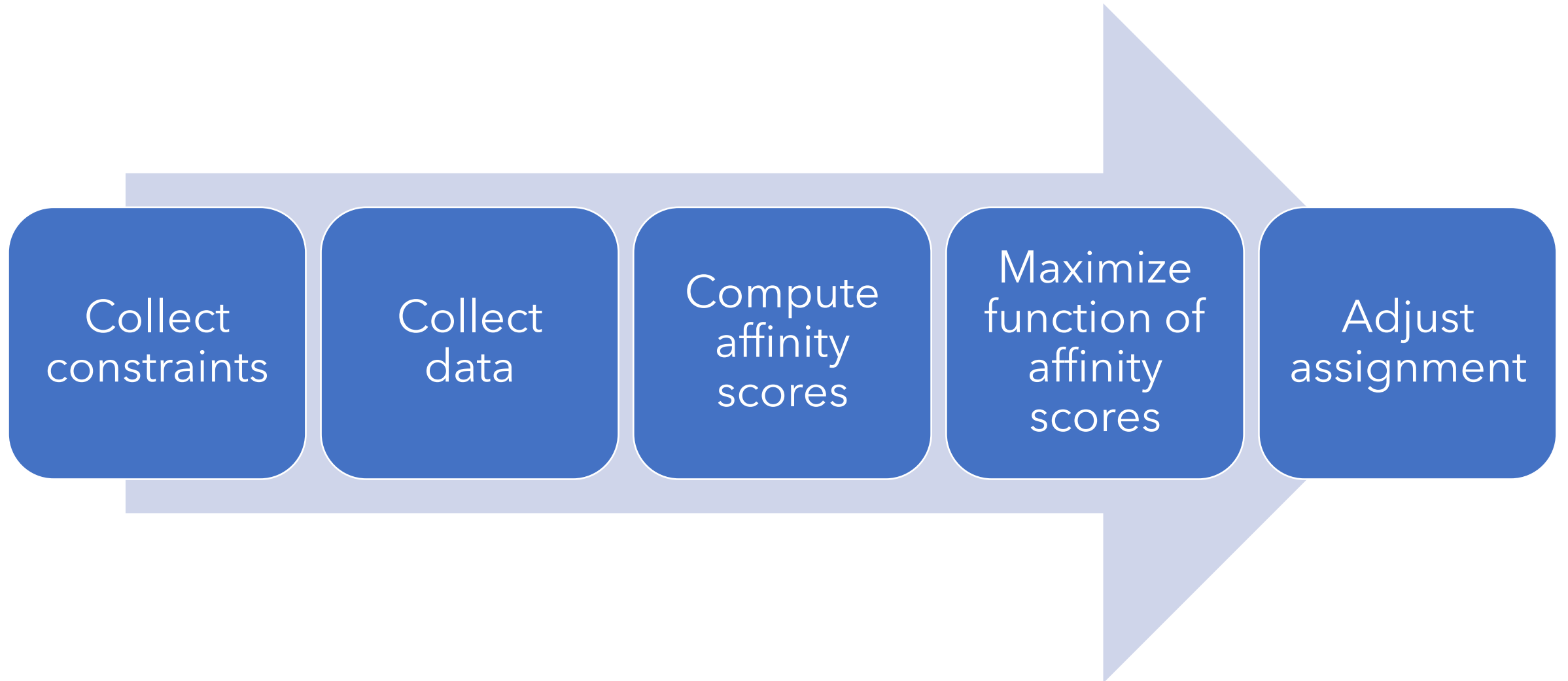
Step through a typical conference workflow



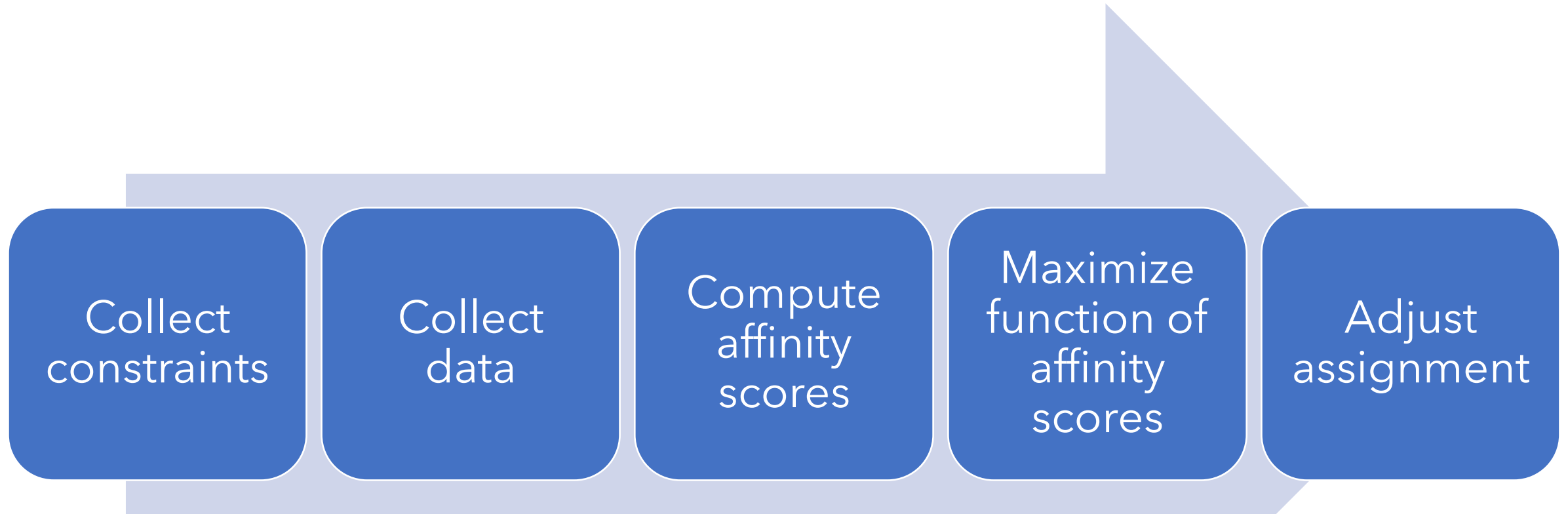
Introduce RRA, a framework that accounts for uncertainty

Also allows us to use new affinity score estimation methods

Conference Workflow 1.0



Conference Workflow 1.0



Kobren et al. Paper matching with local fairness constraints. 2019.

Stelmakh et al. PeerReview4All: Fair and accurate reviewer assignment in peer review. 2019.




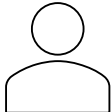



Leyton-Brown et al. Matching papers and reviewers at large conferences. 2022.

Charlin and Zemel. The Toronto paper matching system: an automated paper-reviewer assignment system. 2013



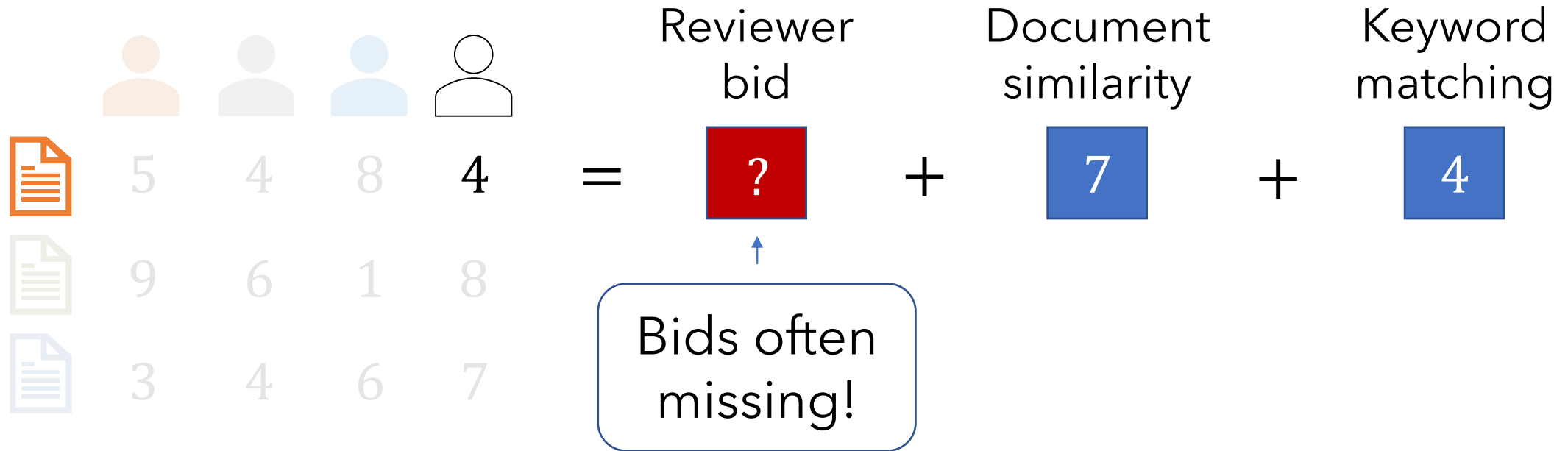


Affinity scores

						Reviewer bid		Document similarity		Keyword matching
	5	4	8	4	=	1	+	7	+	4
	9	6	1	8						
	3	4	6	7						



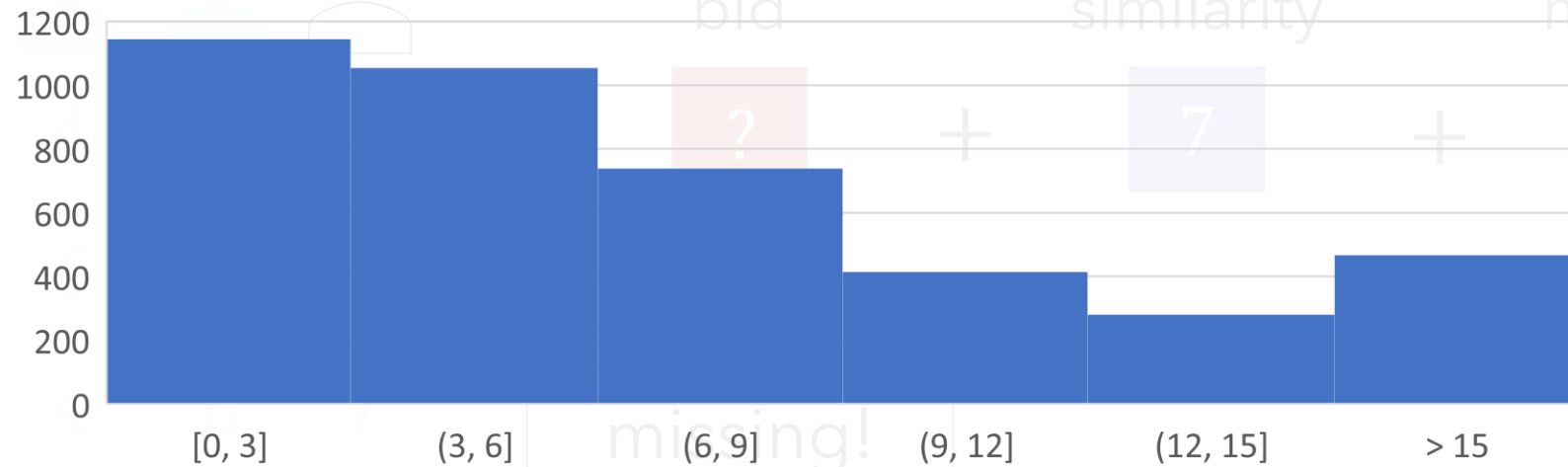
Affinity scores





Affinity scores

Positive Bid Histogram for a Recent Large AI Conference



missing!

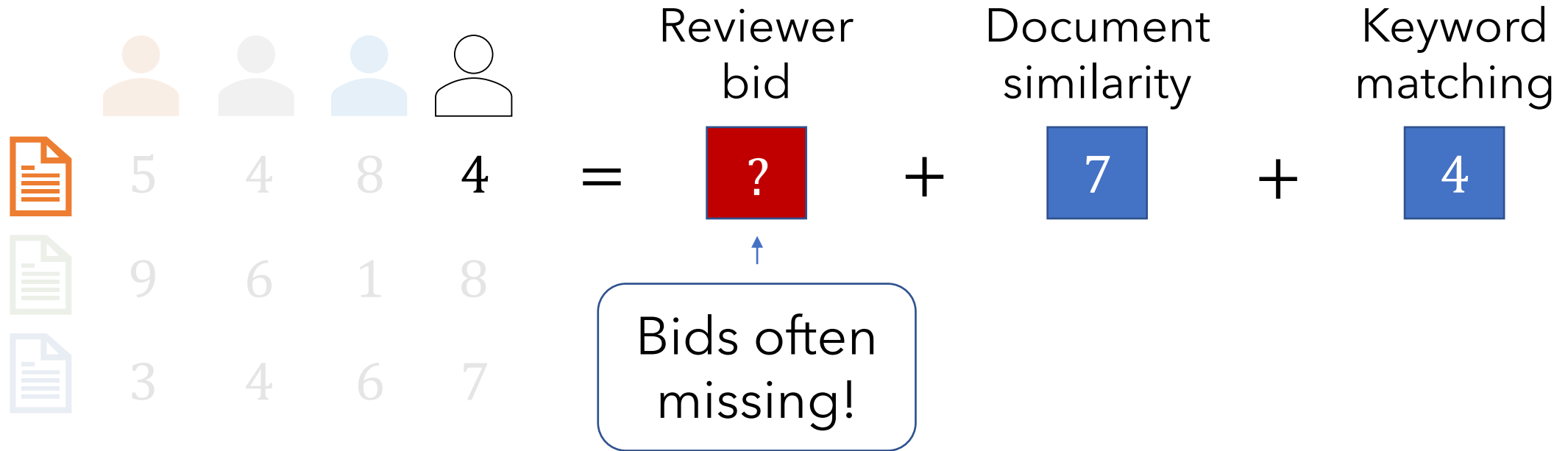
Reviewer bid

Document similarity

Keyword matching

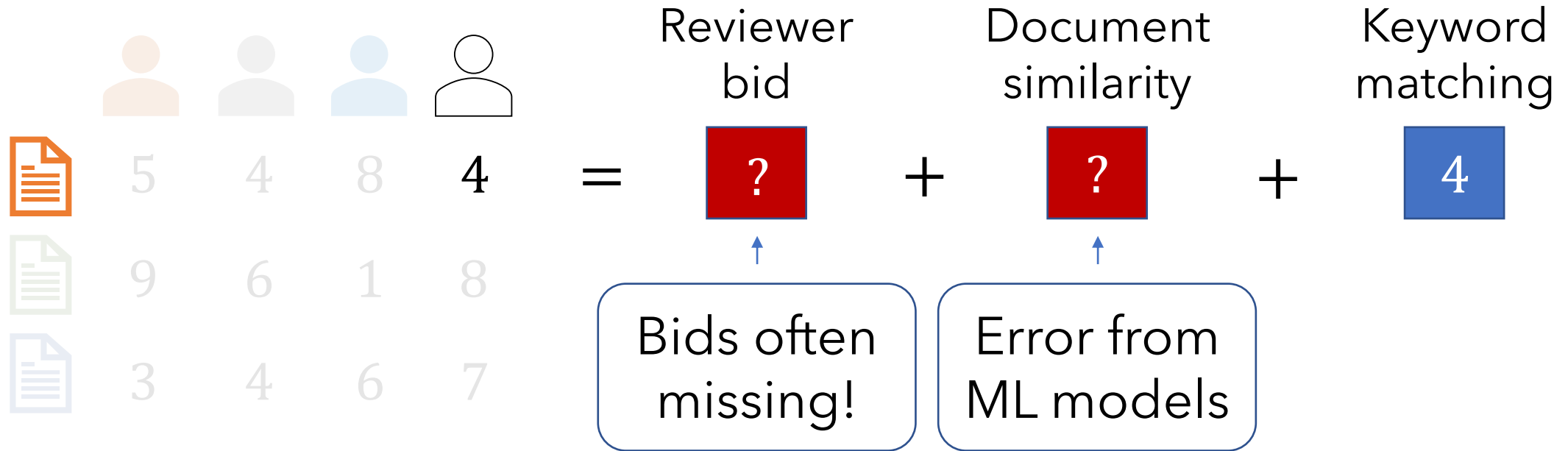


Affinity scores



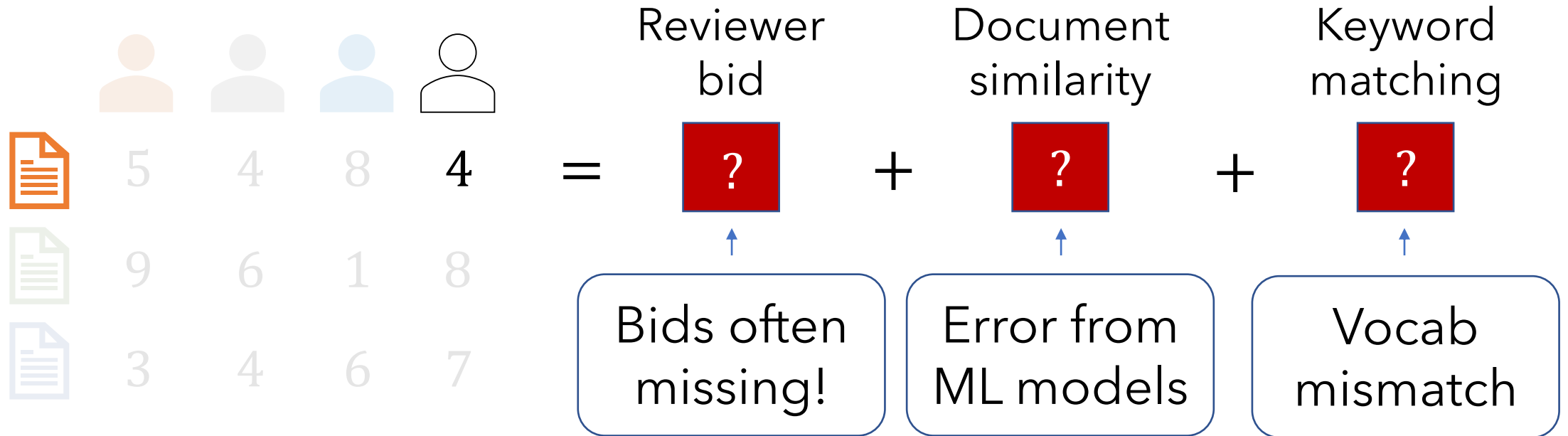


Affinity scores



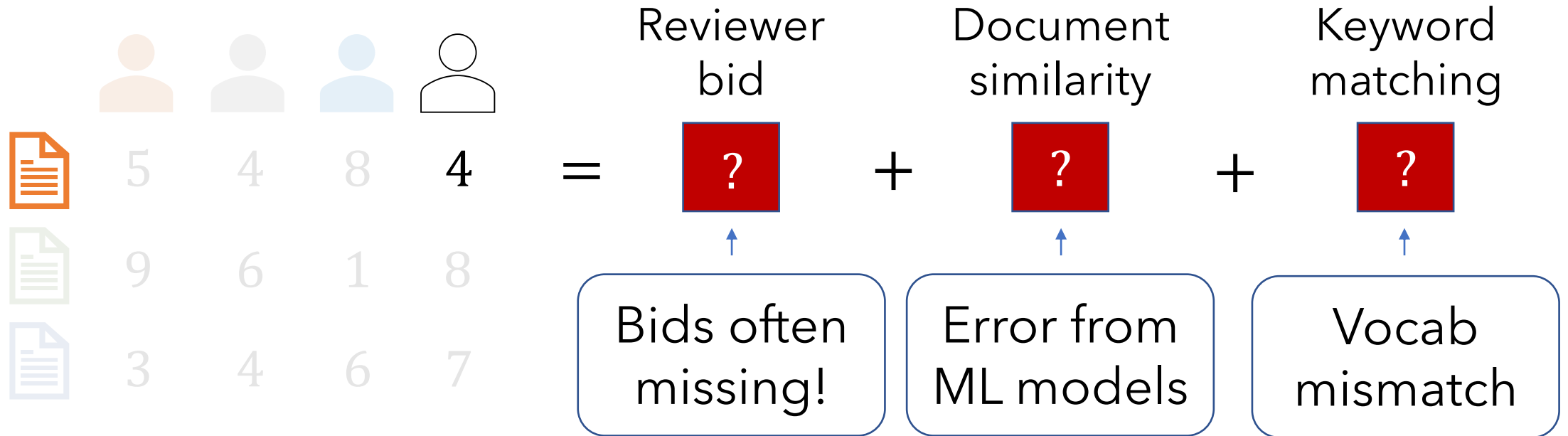


Affinity scores

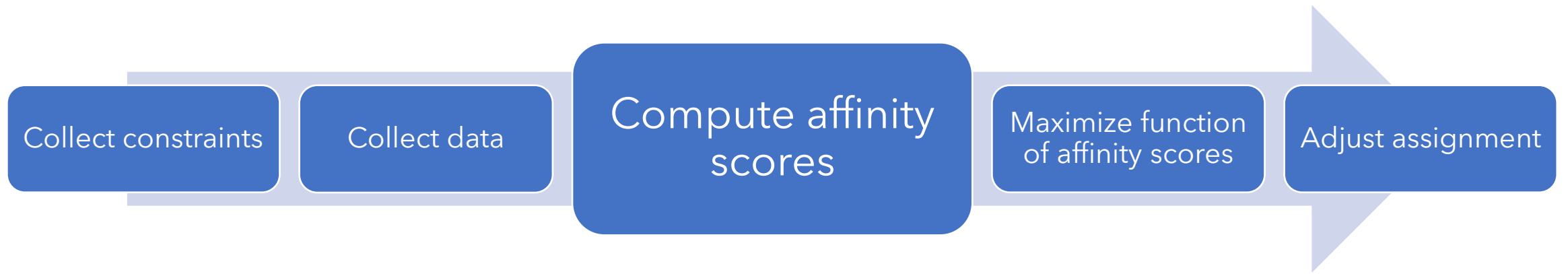


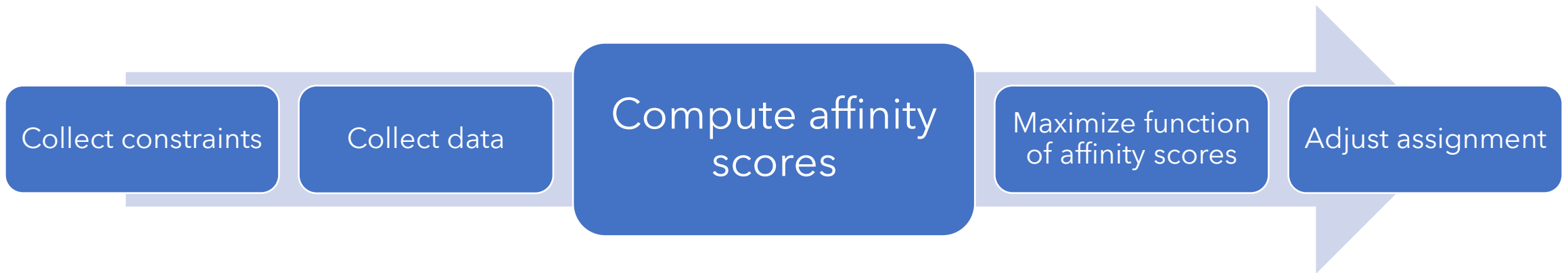


Affinity scores




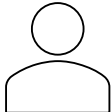





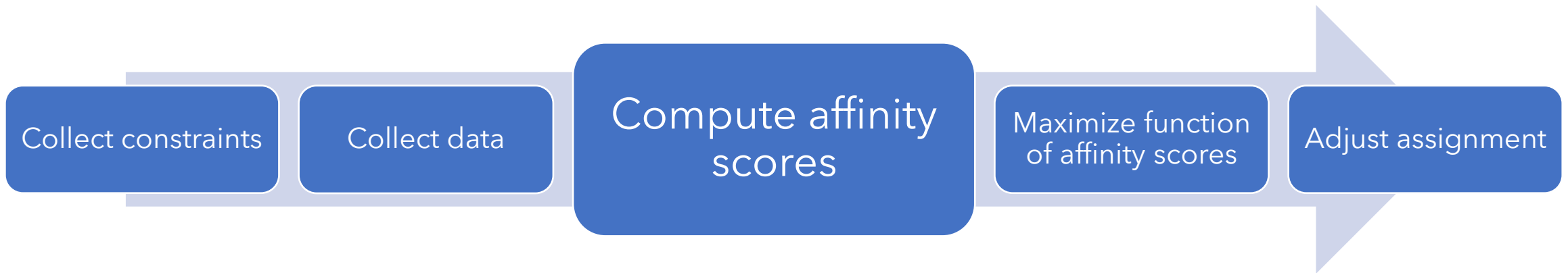
All just proxies!





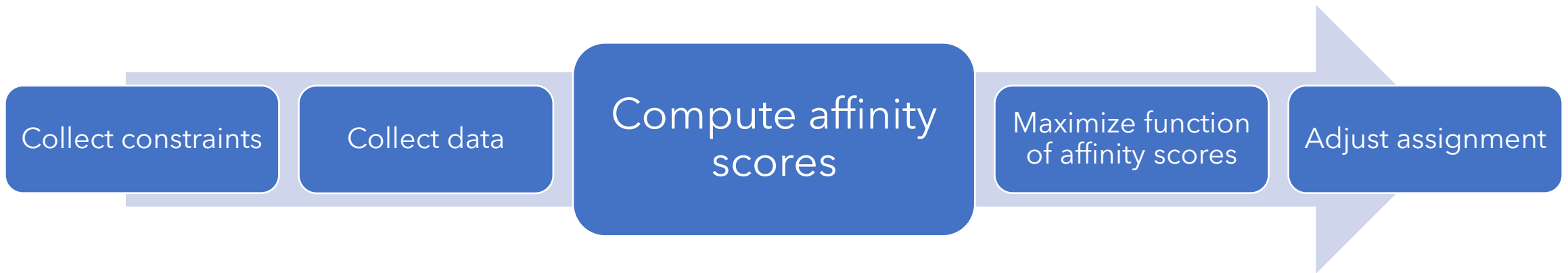
Affinity scores

						Reviewer bid		Document similarity		Keyword matching
	5	4	8	4	=	1	+	7	+	4
	9	6	1	8						
	3	4	6	7						



$$4 = \lambda_1 \begin{matrix} \text{Reviewer} \\ \text{bid} \\ 1 \end{matrix} + \lambda_2 \begin{matrix} \text{Document} \\ \text{similarity} \\ 7 \end{matrix} + \lambda_3 \begin{matrix} \text{Keyword} \\ \text{matching} \\ 4 \end{matrix}$$



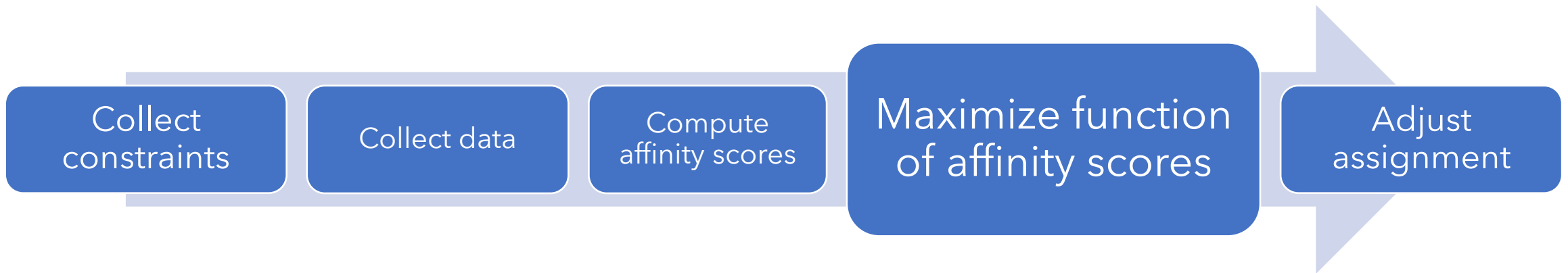


$$\begin{array}{ccccc}
 & \text{Reviewer} & & \text{Document} & & \text{Keyword} \\
 & \text{bid} & & \text{similarity} & & \text{matching} \\
 4 & = \lambda_1 \boxed{1} & + & \lambda_2 \boxed{7} & + & \lambda_3 \boxed{4}
 \end{array}$$




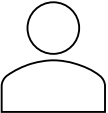





Any unavailable data is set to 0

Underestimation means high affinity scores are high with certainty



Given *affinity scores*

				
	5	4	8	4
	9	6	1	8
	3	4	6	7

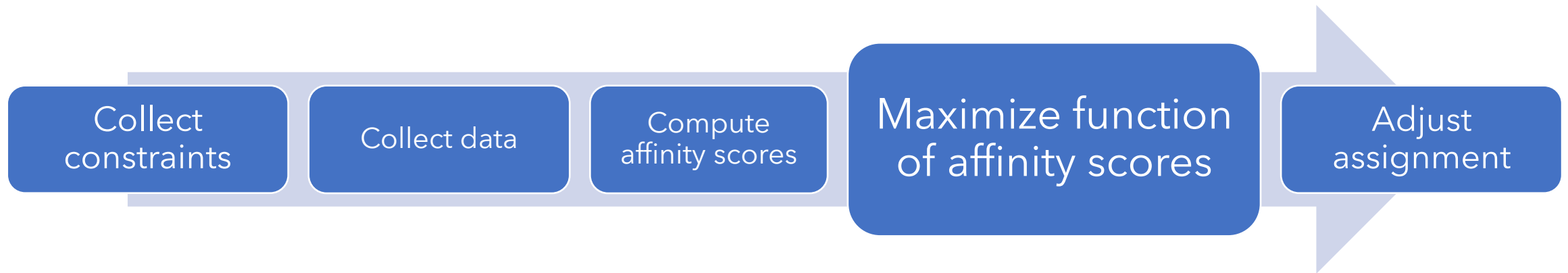
Papers P and reviewers R

Valid assignments $\mathcal{A} \subseteq \{0, 1\}^{|P| \times |R|}$




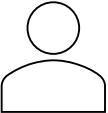



Affinity score matrix $S \in \mathbb{R}_+^{|P| \times |R|}$

Maximize Utilitarian Social Welfare

$$\max_{A \in \mathcal{A}} \left[\sum_{p \in P, r \in R} A_{pr} S_{pr} = USW(A, S) \right]$$



Given *affinity scores*

				
	5	4	8	4
	9	6	1	8
	3	4	6	7

Papers P and reviewers R

Valid assignments $\mathcal{A} \subseteq \{0, 1\}^{|P| \times |R|}$

Affinity score matrix $S \in \mathbb{R}_+^{|P| \times |R|}$

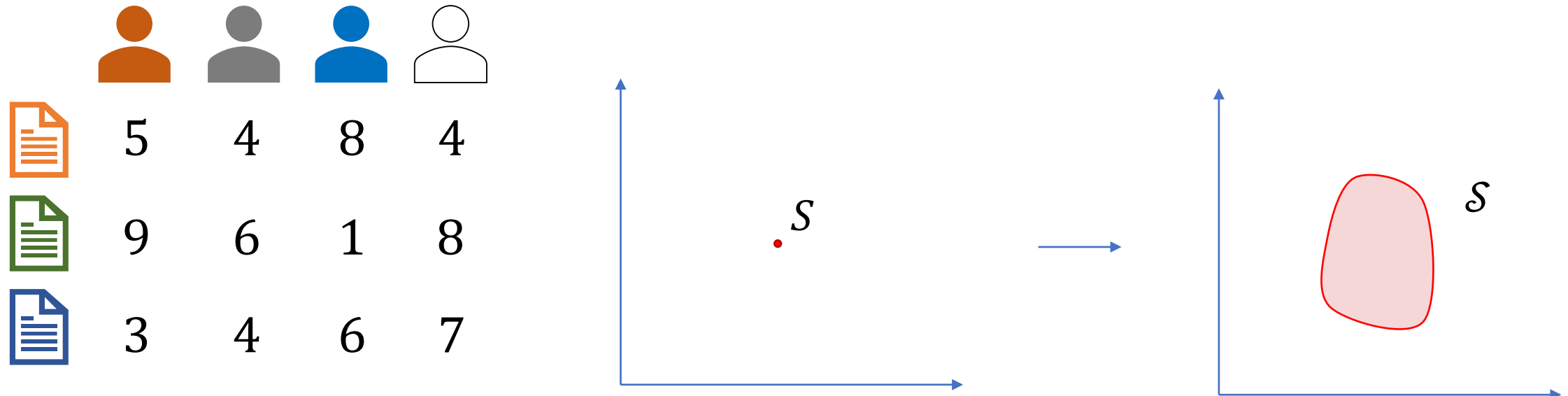
Maximize Utilitarian Social Welfare

$$\max_{A \in \mathcal{A}} \left[\sum_{p \in P, r \in R} A_{pr} S_{pr} = USW(A, S) \right]$$

ILP with totally unimodular constraints
(poly-time solvable)

Robust Reviewer Assignment (RRA)

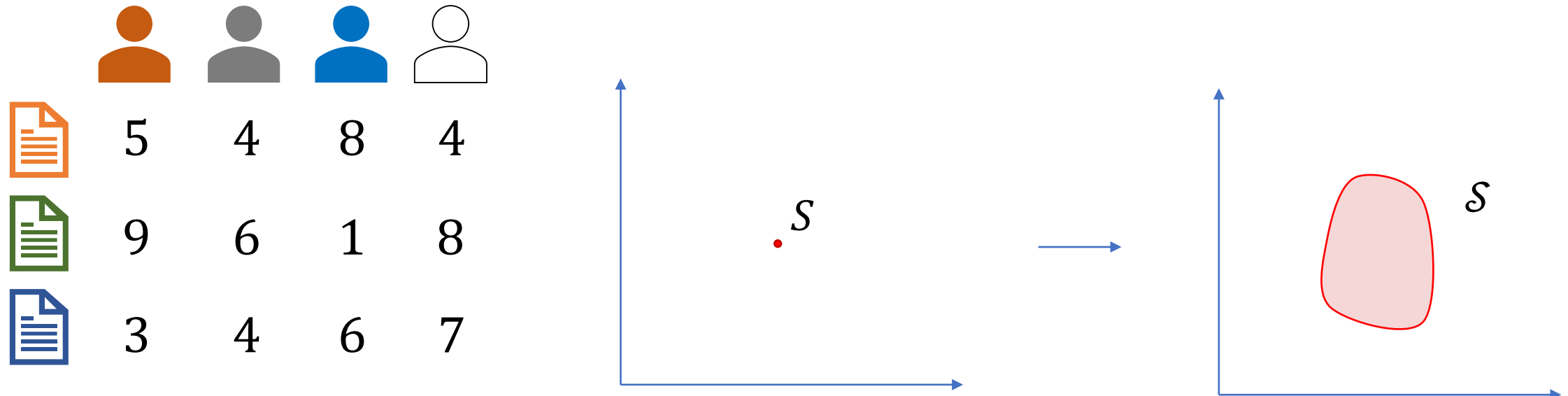
Compute an *uncertainty set* \mathcal{S} containing true, unknown affinity scores S



Maximize the worst-case welfare over the *uncertainty set*

Robust Reviewer Assignment (RRA)

Compute an *uncertainty set* \mathcal{S} containing true, unknown affinity scores S



Maximize the worst-case welfare over the *uncertainty set*

ICLR Experiments

Got all
submitted
papers to last 5
years of ICLR

Used authors
as "reviewers"

Keywords =
paper topics
and author
expertise

International Conference on Learning Representations

ICLR 2021

Oral Presentations

Spotlight Presentations

Poster Presentations

Withdrawn/Rejected Submissions

On the mapping between Hopfield networks and Restricted Boltzmann Mach

Matthew Smart, Anton Zilman

28 Sept 2020 (modified: 10 Feb 2022) ICLR 2021 Oral Readers:  Everyone 9 Replies

Hide details

Keywords: Hopfield Networks, Restricted Boltzmann Machines, Statistical Physics

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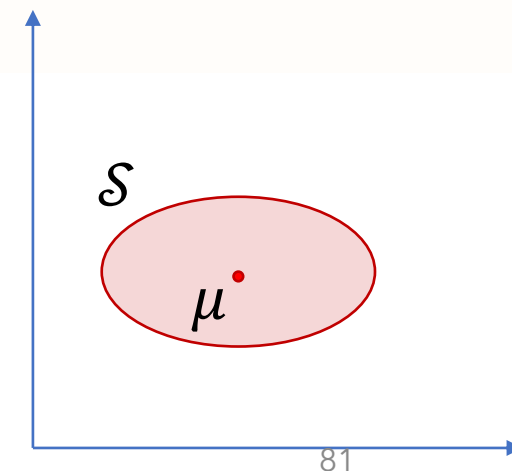
Hide details

Keywords: Hopfield Networks, Restricted Boltzmann Machines, Statistical Physics

Compute \mathcal{S} as a contour of $\mathcal{N}(\mu, \Sigma)$

μ based on keyword overlap

$\text{diag}(\Sigma) \propto 1/(\# \text{ keywords})$



Robust Reviewer Assignment (RRA)

Relax discrete \rightarrow continuous

Subgradient-ascent optimization

Randomized rounding for discrete solution

Robust Reviewer Assignment (RRA)

Relax discrete -> continuous

$$\mathcal{A} \subseteq \{0, 1\}^{m \times n} \rightarrow \tilde{\mathcal{A}} \subseteq [0, 1]^{m \times n}$$

Subgradient-ascent optimization

Randomized rounding for discrete solution

Robust Reviewer Assignment (RRA)

Relax discrete -> continuous

$$\mathcal{A} \subseteq \{0, 1\}^{m \times n} \rightarrow \tilde{\mathcal{A}} \subseteq [0, 1]^{m \times n}$$

Subgradient-ascent optimization

Solve $\max_{A \in \tilde{\mathcal{A}}} f(A)$ by steps in $\partial_A f(A)$

Randomized rounding for discrete solution

Robust Reviewer Assignment (RRA)

Subgradient-ascent optimization

Solve $\max_{A \in \tilde{\mathcal{A}}} f(A)$ by steps in $\partial_A f(A)$

Robust Reviewer Assignment (RRA)

Subgradient-ascent optimization

Solve $\max_{A \in \tilde{\mathcal{A}}} f(A)$ by steps in $\partial_A f(A)$

Robust Reviewer Assignment (RRA)

Subgradient-ascent optimization

T iterations

Compute $f(A)$ as the adversary

Step in direction of $\partial_A f(A)$

Project to feasible set $\tilde{\mathcal{A}}$

Solve $\max_{A \in \tilde{\mathcal{A}}} f(A)$ by steps in $\partial_A f(A)$

Robust Reviewer Assignment (RRA)

Subgradient-ascent optimization

T iterations

Compute $f(A)$ as the adversary

Step in direction of $\partial_A f(A)$

Project to feasible set $\tilde{\mathcal{A}}$

Solve $\max_{A \in \tilde{\mathcal{A}}} f(A)$ by steps in $\partial_A f(A)$

$$f(A) = \min_{S \in \mathcal{S}} USW(A, S)$$

Robust Reviewer Assignment (RRA)

Subgradient-ascent optimization

T iterations

Compute $f(A)$ as the adversary

Step in direction of $\partial_A f(A)$

Project to feasible set $\tilde{\mathcal{A}}$

Solve $\max_{A \in \tilde{\mathcal{A}}} f(A)$ by steps in $\partial_A f(A)$

$$f(A) = \min_{S \in \mathcal{S}} USW(A, S)$$

$$\nabla_A USW(A, \operatorname{argmin}_{S \in \mathcal{S}} USW(A, S)) \in \partial_A f(A)$$

$$A' \leftarrow A + \nabla_A USW(A, \operatorname{argmin}_{S \in \mathcal{S}} USW(A, S))$$

Robust Reviewer Assignment (RRA)

Subgradient-ascent optimization

T iterations

Compute $f(A)$ as the adversary

Step in direction of $\partial_A f(A)$

Project to feasible set $\tilde{\mathcal{A}}$

Solve $\max_{A \in \tilde{\mathcal{A}}} f(A)$ by steps in $\partial_A f(A)$

$$f(A) = \min_{S \in \mathcal{S}} USW(A, S)$$

$$\nabla_A USW(A, \operatorname{argmin}_{S \in \mathcal{S}} USW(A, S)) \in \partial_A f(A)$$

$$A' \leftarrow A + \nabla_A USW(A, \operatorname{argmin}_{S \in \mathcal{S}} USW(A, S))$$

$$A'' \leftarrow \operatorname{argmin}_{A \in \tilde{\mathcal{A}}} \|A - A'\|_2$$

Robust Reviewer Assignment (RRA)

Relax discrete \rightarrow continuous

$$\mathcal{A} \subseteq \{0, 1\}^{m \times n} \rightarrow \tilde{\mathcal{A}} \subseteq [0, 1]^{m \times n}$$

Subgradient-ascent optimization

Solve $\max_{A \in \tilde{\mathcal{A}}} f(A)$ by steps in $\partial_A f(A)$

Randomized rounding for discrete solution

Round $A \in \tilde{\mathcal{A}}$ to $A' \in \mathcal{A}$

Spherical Uncertainty Sets

Theorem 2:

$$\max_{A \in \mathcal{A}} \min_{S \in \mathcal{B}_\epsilon(S_0)} USW(A, S) = \max_{A \in \mathcal{A}} USW(A, S_0)$$

Proof sketch:

$$\min_{S \in \mathcal{B}_\epsilon(S_0)} USW(A, S) = \min_{X \in \mathcal{B}_\epsilon(0)} USW(A, S_0 + X)$$

Spherical Uncertainty Sets

Theorem 2:

$$\max_{A \in \mathcal{A}} \min_{S \in \mathcal{B}_\epsilon(S_0)} USW(A, S) = \max_{A \in \mathcal{A}} USW(A, S_0)$$

Proof sketch:

$$\begin{aligned} \min_{S \in \mathcal{B}_\epsilon(S_0)} USW(A, S) &= \min_{X \in \mathcal{B}_\epsilon(0)} USW(A, S_0 + X) \\ &= USW(A, S_0) + \min_{X \in \mathcal{B}_\epsilon(0)} USW(A, X) \end{aligned}$$

Spherical Uncertainty Sets

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$$\max_{A \in \mathcal{A}} \min_{S \in \mathcal{B}_\epsilon(S_0)} USW(A, S) = \max_{A \in \mathcal{A}} USW(A, S_0)$$

Proof sketch:

$$\min_{S \in \mathcal{B}_\epsilon(S_0)} USW(A, S) = \min_{X \in \mathcal{B}_\epsilon(0)} USW(A, S_0 + X)$$

Minimized when X points
as far in the direction
opposite A as possible:

$$X = -A \frac{\epsilon}{\|A\|_2}$$

$$= USW(A, S_0) + \min_{X \in \mathcal{B}_\epsilon(0)} USW(A, X)$$


Spherical Uncertainty Sets

Theorem 2:

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$$= USW(A, S_0) + USW\left(A, -A \frac{\epsilon}{\|A\|_2}\right)$$

Spherical Uncertainty Sets

Theorem 2:

$$\max_{A \in \mathcal{A}} \min_{S \in \mathcal{B}_\epsilon(S_0)} USW(A, S) = \max_{A \in \mathcal{A}} USW(A, S_0)$$

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$$= USW(A, S_0) + USW\left(A, -A \frac{\epsilon}{\|A\|_2}\right)$$

$$= USW(A, S_0) - \|A\|_2^2 \frac{\epsilon}{\|A\|_2}$$

Spherical Uncertainty Sets

Theorem 2:

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$$= USW(A, S_0) - \|A\|_2 \epsilon$$

Spherical Uncertainty Sets

Theorem 2:

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Proof sketch:

$$\min_{S \in \mathcal{B}_\epsilon(S_0)} USW(A, S) = \min_{X \in \mathcal{B}_\epsilon(0)} USW(A, S_0 + X)$$

Minimized when X points
as far in the direction
opposite A as possible:

$$X = -A \frac{\epsilon}{\|A\|_2}$$

Constant, due to
constraints

$$= USW(A, S_0) + \min_{X \in \mathcal{B}_\epsilon(0)} USW(A, X)$$

$$= USW(A, S_0) + USW\left(A, -A \frac{\epsilon}{\|A\|_2}\right)$$

$$= USW(A, S_0) - \|A\|_2^2 \frac{\epsilon}{\|A\|_2}$$

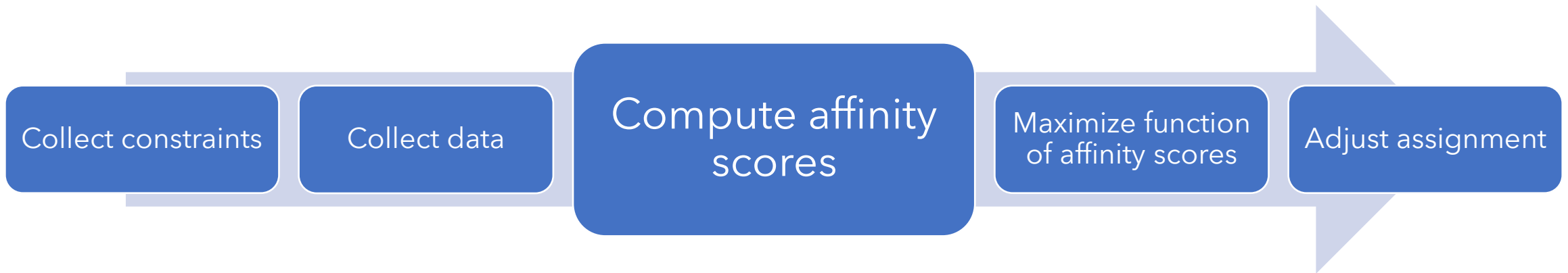
$$= USW(A, S_0) - \|A\|_2 \epsilon$$

Future Work

- Be sure to publicize FairSeq, our recently submitted paper, and the workshop.
- Further experiments
 - Take datasets of affinity scores, drop some and fill in remaining values
 - Build a predictive model of affinity scores from old conferences
 - Collect data from conference organizers
 - Discuss the collaborative filtering bid model?
- Combining fairness and robustness
 - Add an envy penalty to the objective function
 - Maximize expected egalitarian welfare
- Rounding tradeoffs?

Fair Division at IJCAI

- Anyone want to co-organize?
- Plan to submit papers!
- Mention proposed schedule, etc.



Affinity scores

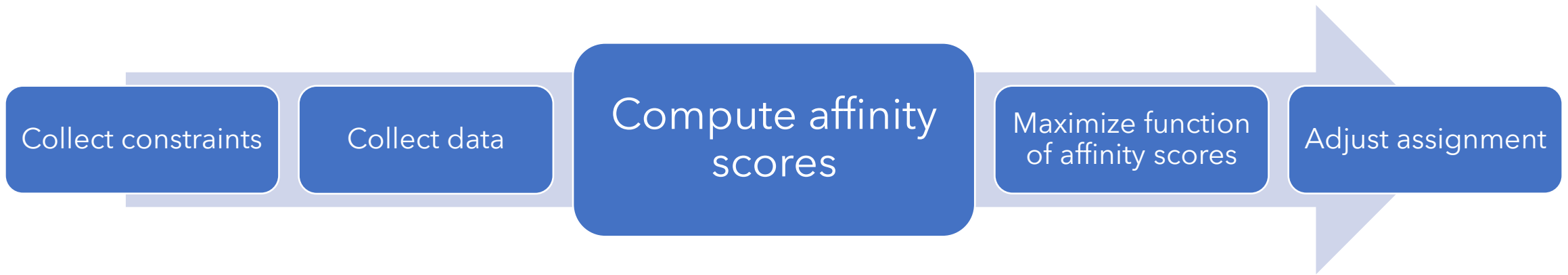


	Reviewer bid		Document similarity		Keyword matching
4	$= \lambda_1$	1	+	λ_2	7
			+	λ_3	4



$$\text{Base Aggregated Score} = \begin{cases} \frac{1}{4}\text{TPMS} + \frac{1}{4}\text{ACL} + \frac{1}{2}\text{SAM} & \text{if all data is present} \\ \frac{1}{2}\text{ACL} + \frac{1}{2}\text{SAM} & \text{if TPMS is missing} \\ \frac{1}{2}\text{TPMS} + \frac{1}{2}\text{SAM} & \text{if ACL is missing} \\ \text{SAM} & \text{if both ACL and TPMS are missing} \end{cases}$$

$$\text{aggscore} = (\text{Base aggregated score})^{\text{bidscore}}$$



Affinity scores



Reviewer bid Document similarity Keyword matching

4 = λ_1

Does this necessarily correlate with review quality?

Base Aggregated Score

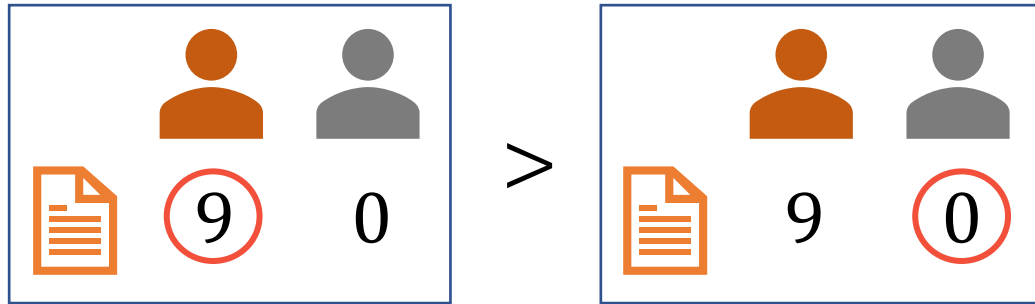
(SAM

If both ACL and TPMS are missing

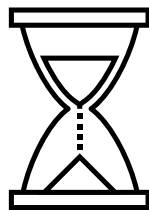
$$\text{aggscore} = (\text{Base aggregated score})^{\text{bidscore}}$$

Allocation Deciderata

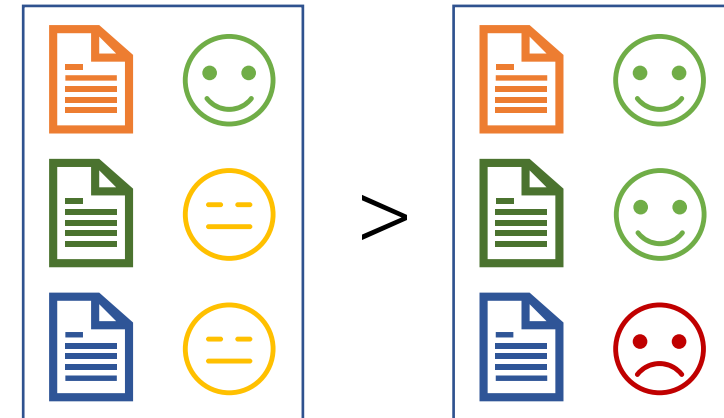
High affinity (welfare/USW)



Fast to compute



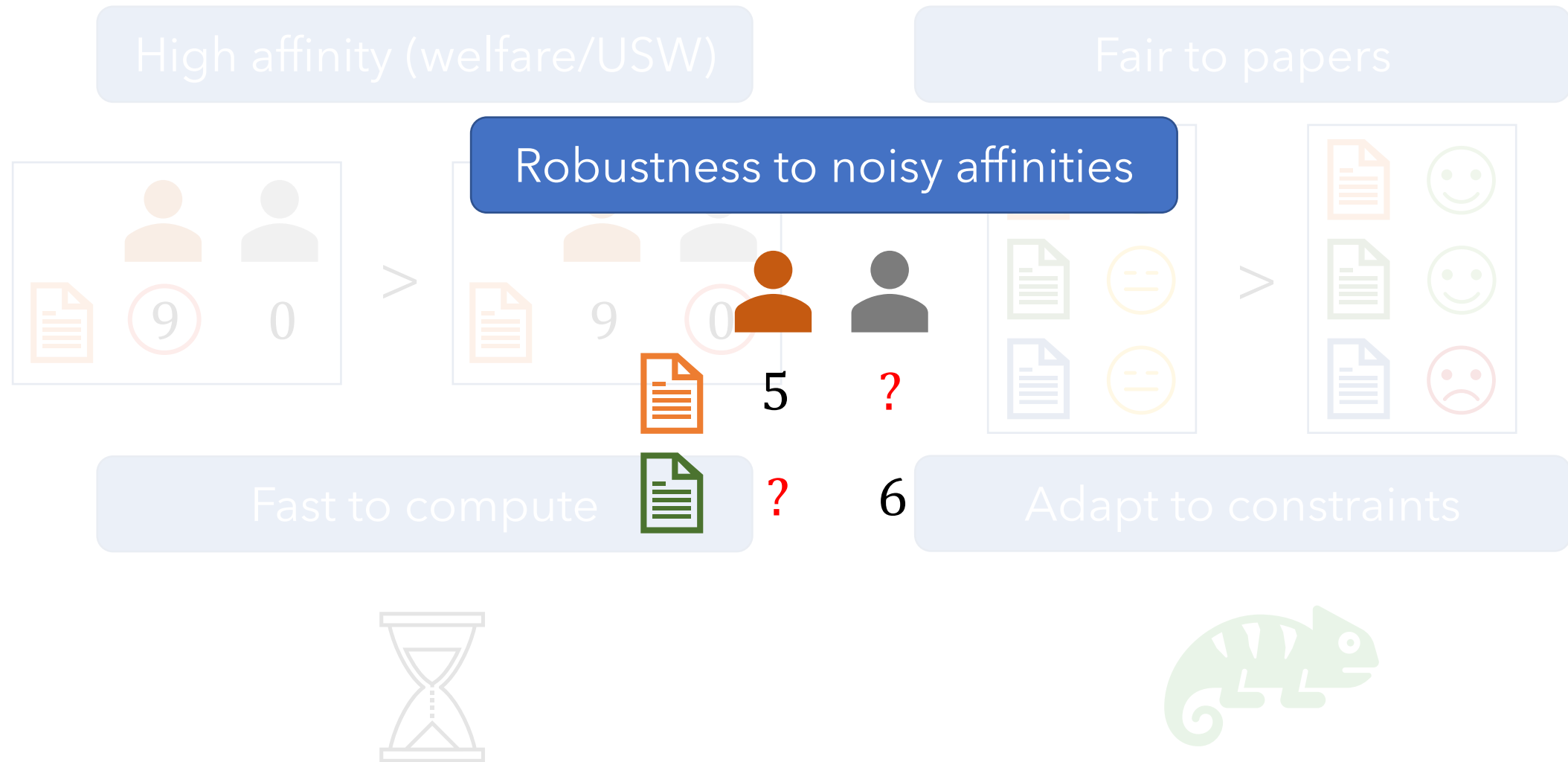
Fair to papers



Adapt to constraints



Allocation Deciderata



Fairness to Papers

Envy-freeness up to 1 Item (EF1)

Ensures pairwise balance in affinity scores




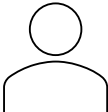


Large literature with simple/fast algorithms

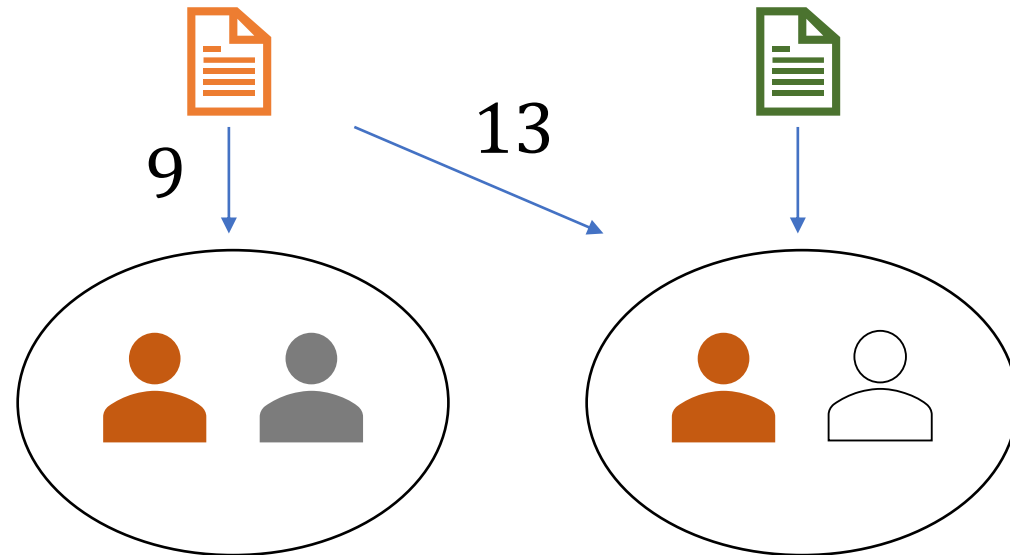
Envy-freeness up to 1 Item (EF1)

For all pairs of papers A and B, if B receives a set of reviewers that is better suited for A, then it is due to at most one reviewer

Envy-freeness up to 1 Item (EF1)




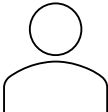


For all pairs of papers A and B, if B receives a set of reviewers that is better suited for A, then it is due to at most one reviewer

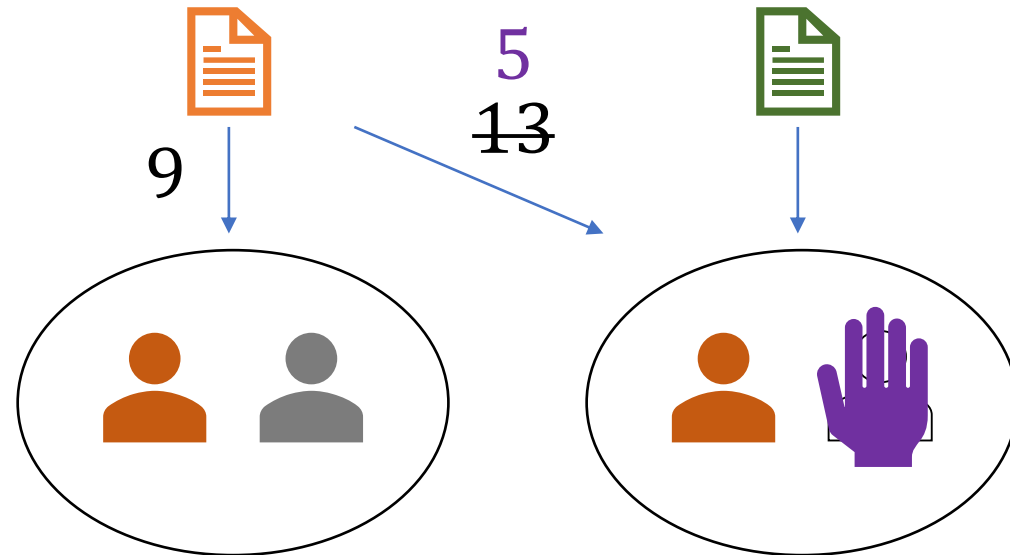
				
	5	4	8	8
	9	6	1	8



Envy-freeness up to 1 Item (EF1)

For all pairs of papers A and B, if B receives a set of reviewers that is better suited for A, then it is due to at most one reviewer

				
	5	4	8	8
	9	6	1	8



Envy-freeness up to 1 Item (EF1)

For all pairs of papers A and B, if B receives a set of reviewers that is better suited for A, then it is due to at most one reviewer

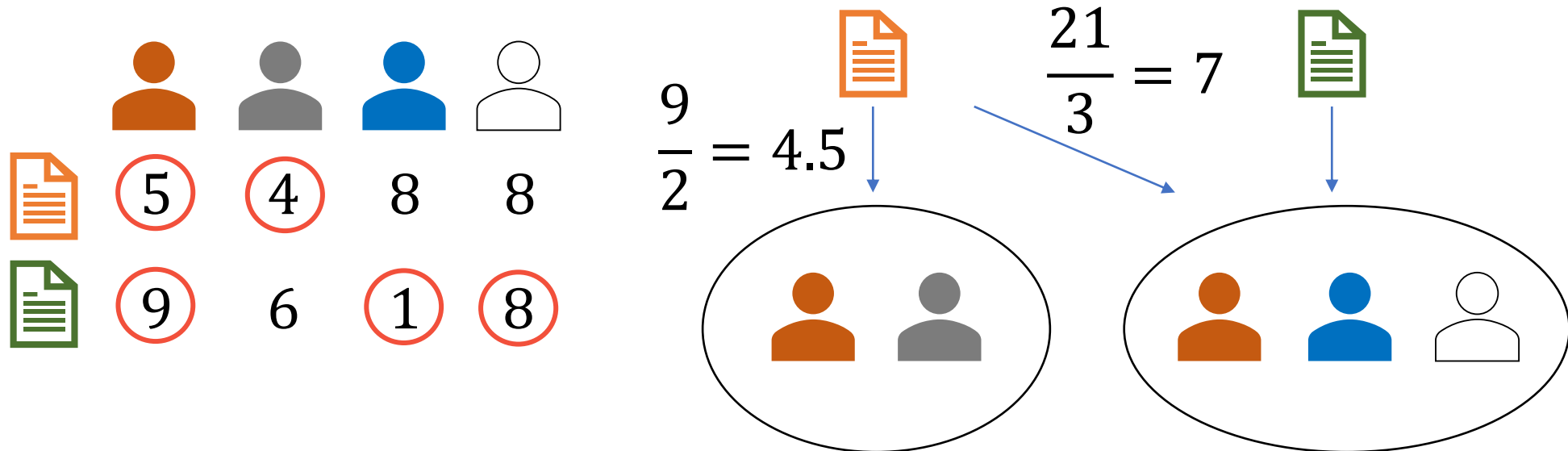
Does not make sense when
papers have variable demands

Weighted EF1 (WEF1)

For all pairs of papers A and B, if B receives a set of reviewers that is better suited for A **after adjusting for paper demands**, then it is due to at most one reviewer

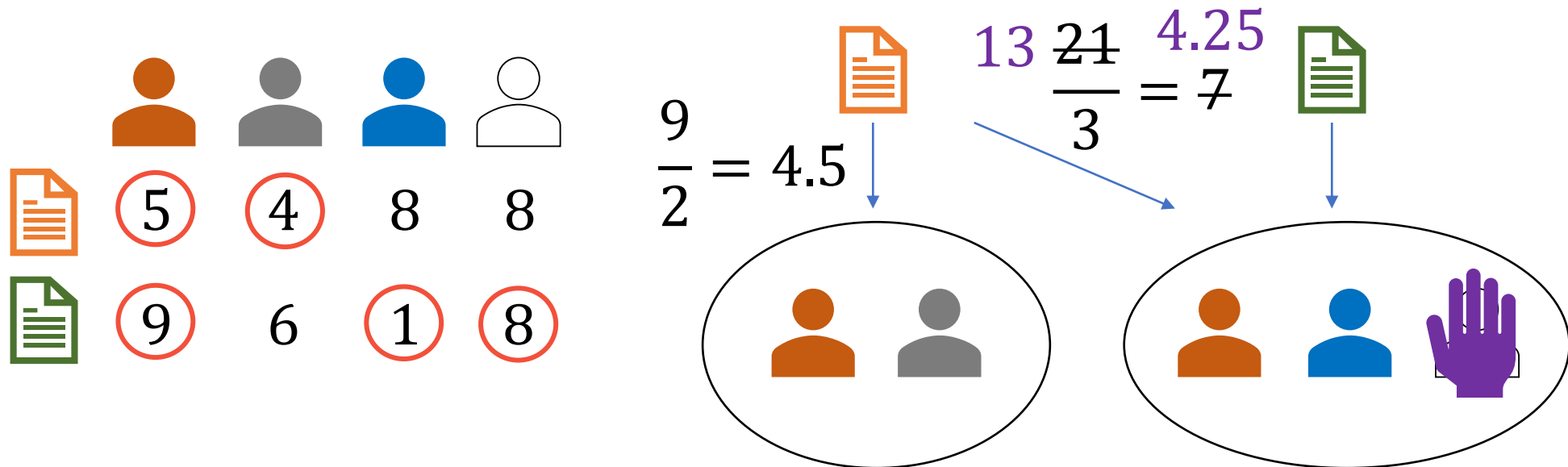
Weighted EF1 (WEF1)

For all pairs of papers A and B, if B receives a set of reviewers that is better suited for A **after adjusting for paper demands**, then it is due to at most one reviewer



Weighted EF1 (WEF1)

For all pairs of papers A and B, if B receives a set of reviewers that is better suited for A **after adjusting for paper demands**, then it is due to at most one reviewer



Picking Sequences for EF1 and Welfare

Put papers in some order,
and assign reviewers in that order

Choose the order for EF1 &
(approximate) max welfare

Picking Sequences for EF1 and Welfare

I Will Have Order! Optimizing Orders for Fair Reviewer Assignment.
Payan and Zick, IJCAI 2022.

Fixed order, repeated over rounds (*round-robin order*) is EF1



Does not work for
variable paper demands

Finding a high-welfare
round-robin order is slow

FairSequence




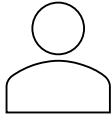



Picking sequence that assigns in order of fraction of demand satisfied

Ties broken to greedily maximize affinity

FairSequence

Picking sequence that assigns in order of fraction of demand satisfied




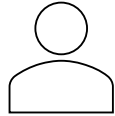



Ties broken to greedily maximize affinity

					
0/2		5	4	8	4
0/2		9	6	1	8
0/3		3	4	6	7

FairSequence

Picking sequence that assigns in order of fraction of demand satisfied




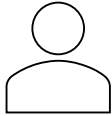



Ties broken to greedily maximize affinity

					
0/2		5	4	8	4
1/2		9	6	1	8
0/3		3	4	6	7

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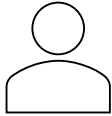



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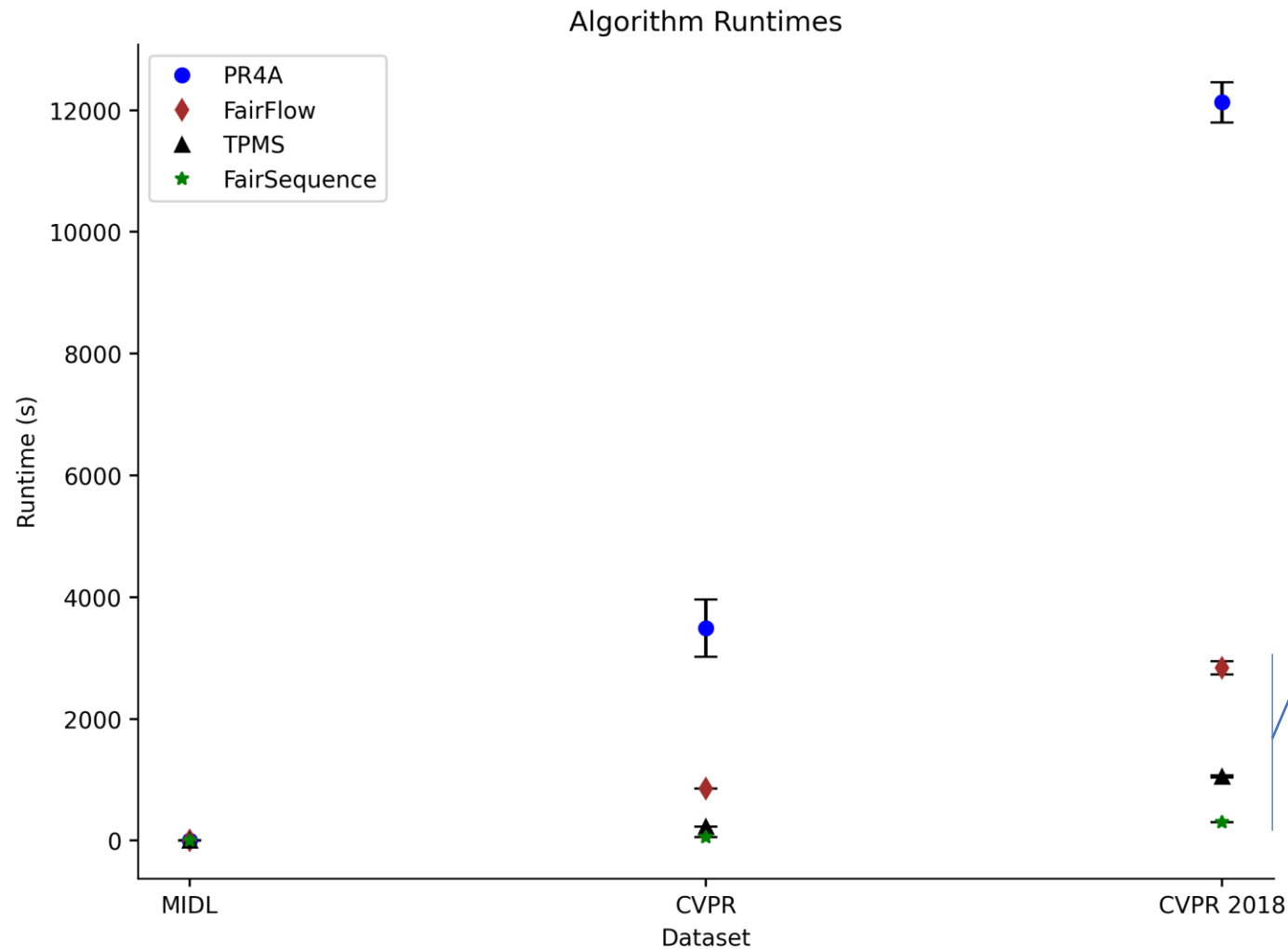
Ties broken to greedily maximize affinity

Very fast

Satisfies WEF1

High welfare in practice

FairSequence – Very Fast!



~10x faster than
FairFlow!

~3x faster than TPMS!

Welfare and Fairness

Our Approaches

		TPMS (OPT)	FairFlow	PR4A	GRRR	FairSeq
MIDL	USW (% of OPT)	100%	100%	98%	98%	99%
	# EF1 Viol.	0	0	0	0	0
		TPMS (OPT)	FairFlow	PR4A	GRRR	FairSeq
CVPR	USW (% of OPT)	100%	96%	94%	88%	92%
	# EF1 Viol.	473545	23344	82	0	0
		TPMS (OPT)	FairFlow	PR4A	GRRR	FairSeq
CVPR '18	USW (% of OPT)	100%	97%	97%	94%	96%
	# EF1 Viol.	134	25	2	0	0

FairSequence Fits the Criteria

Fairness

Ours are the only approaches that satisfy EF1

Welfare

High USW w.r.t. TPMS (OPT) and algorithms used in practice

Speed

~10x speedup compared to fair competitors

Flexibility

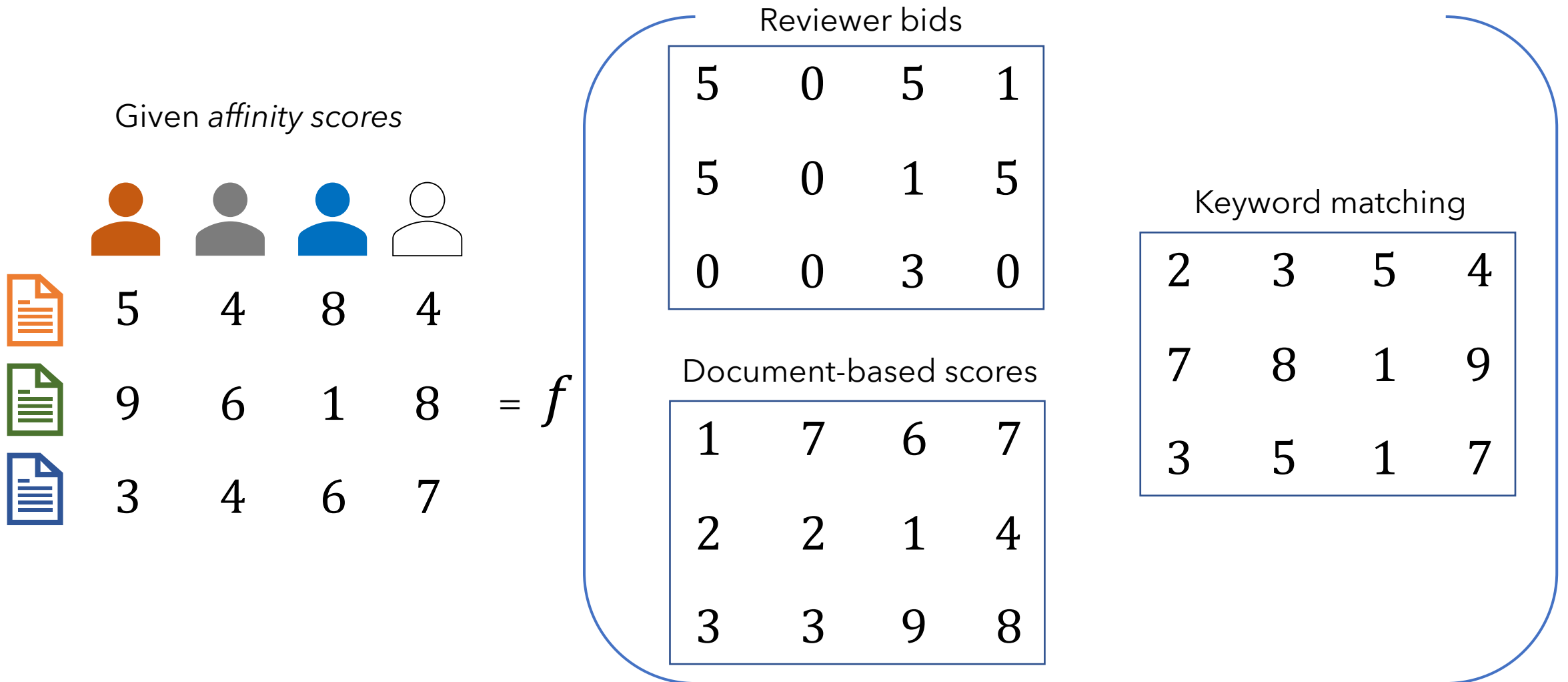
Simplicity → flexibility

Look for FairSequence!

Available in



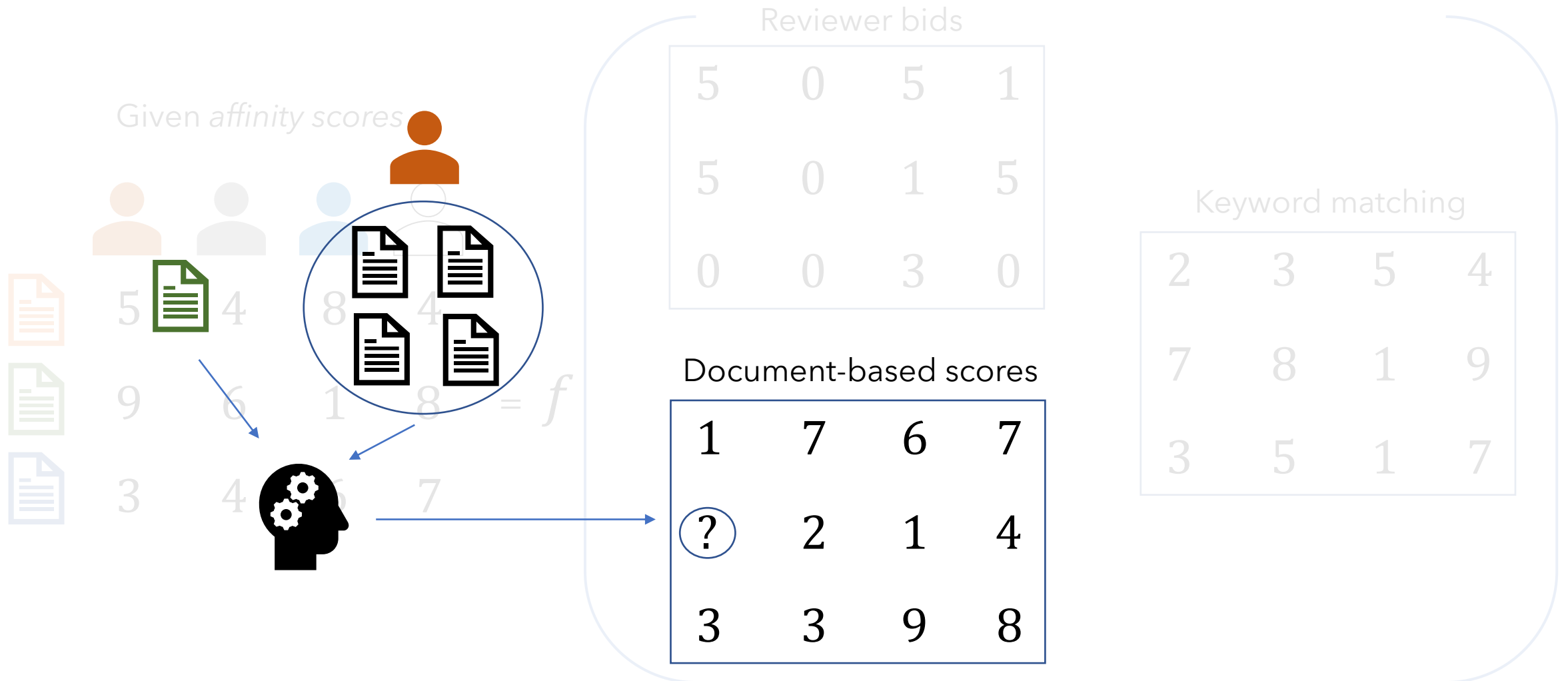
Can we trust affinity scores?



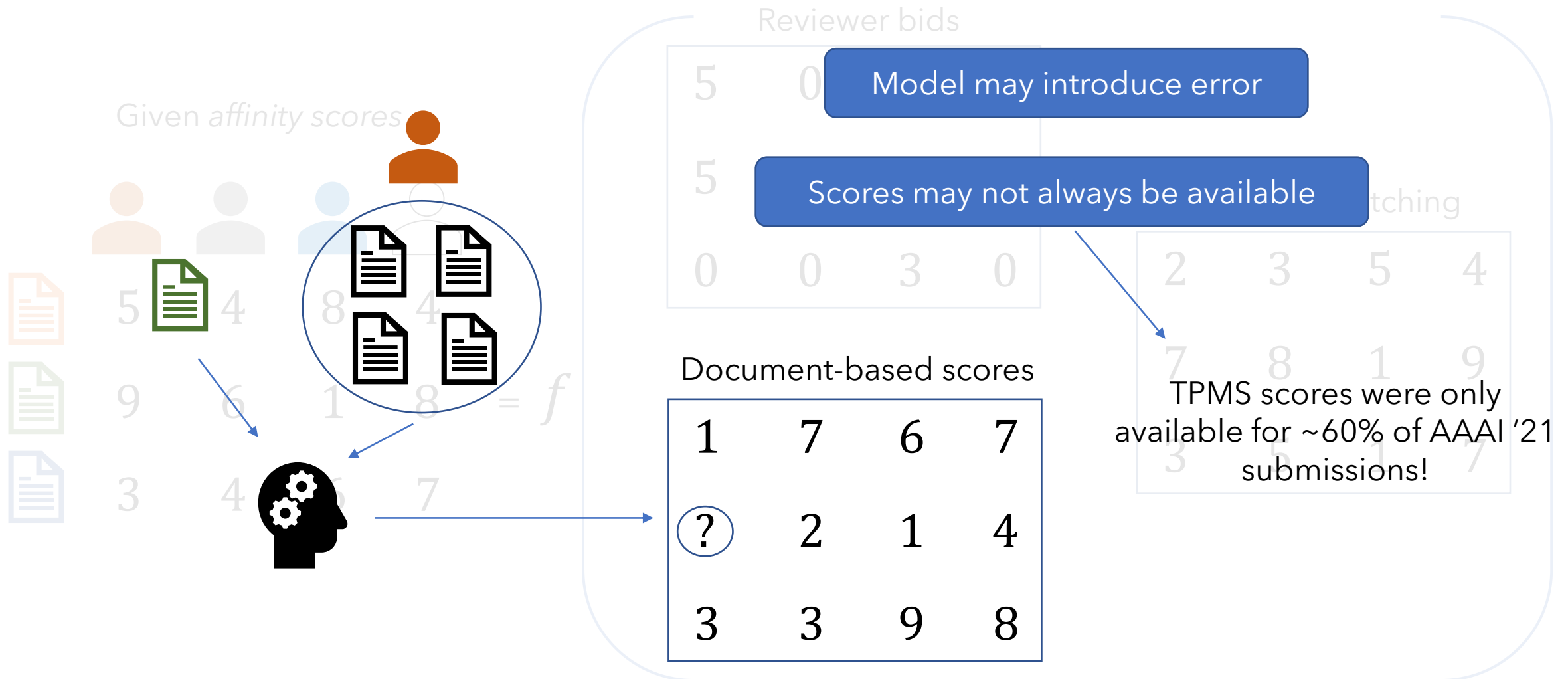
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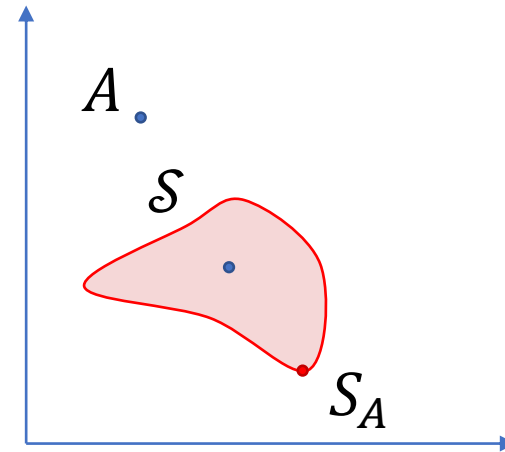
Can we trust affinity scores?



Robust Reviewer Assignment

Feasible region \mathcal{S} for affinity scores
Solution space \mathcal{A} of valid assignments

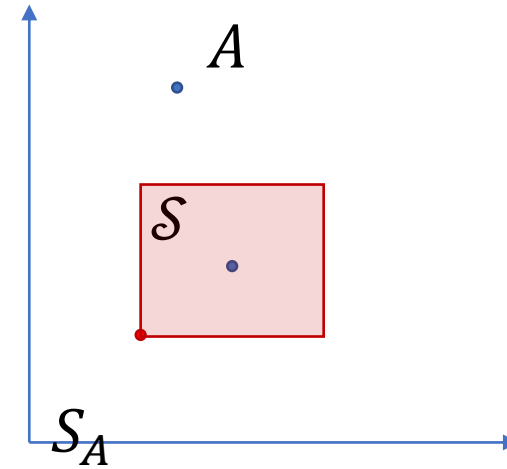
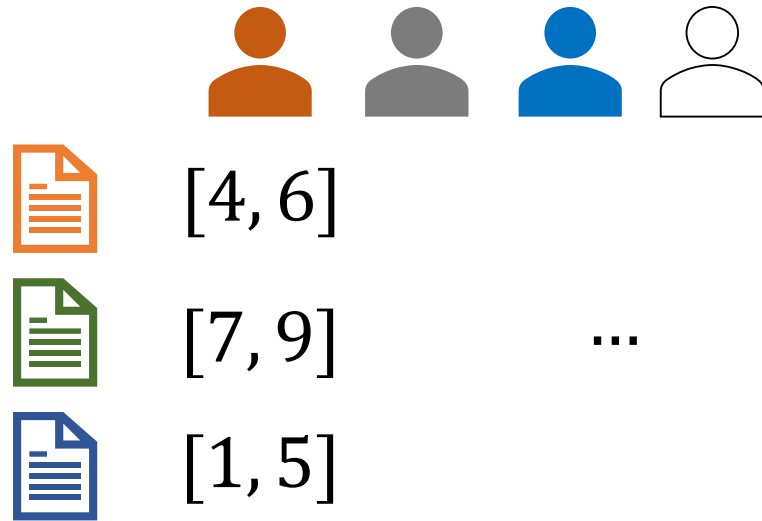
$$\max_{A \in \mathcal{A}} \min_{S_A \in \mathcal{S}} USW(A, S_A)$$



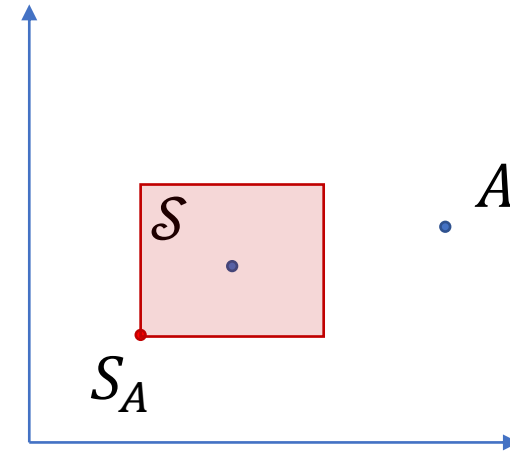
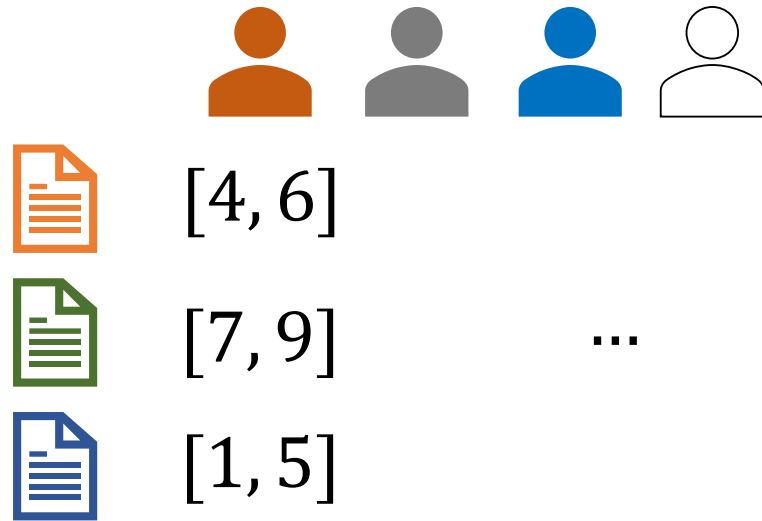
Goals:

1. Find plausible models for \mathcal{S} with simple solutions
2. Build general purpose robust optimization algorithm
3. Enable alternative welfare functions

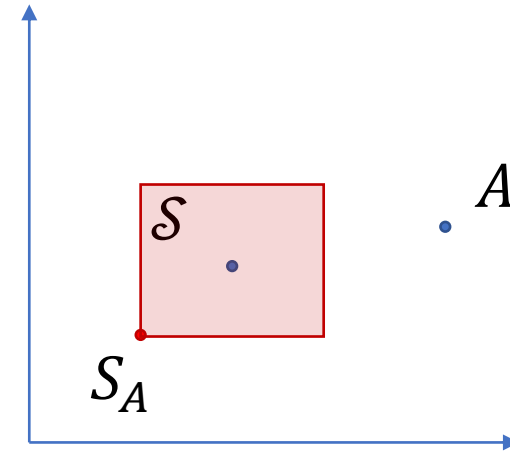
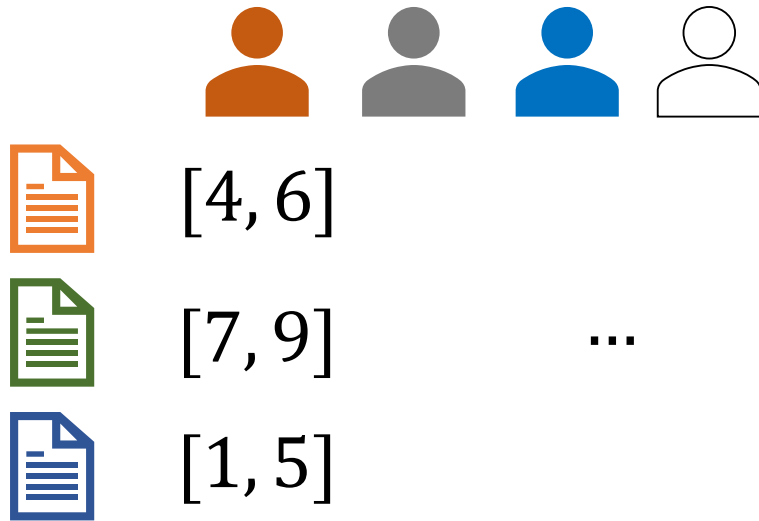
\mathcal{S} as Box Constraints



\mathcal{S} as Box Constraints



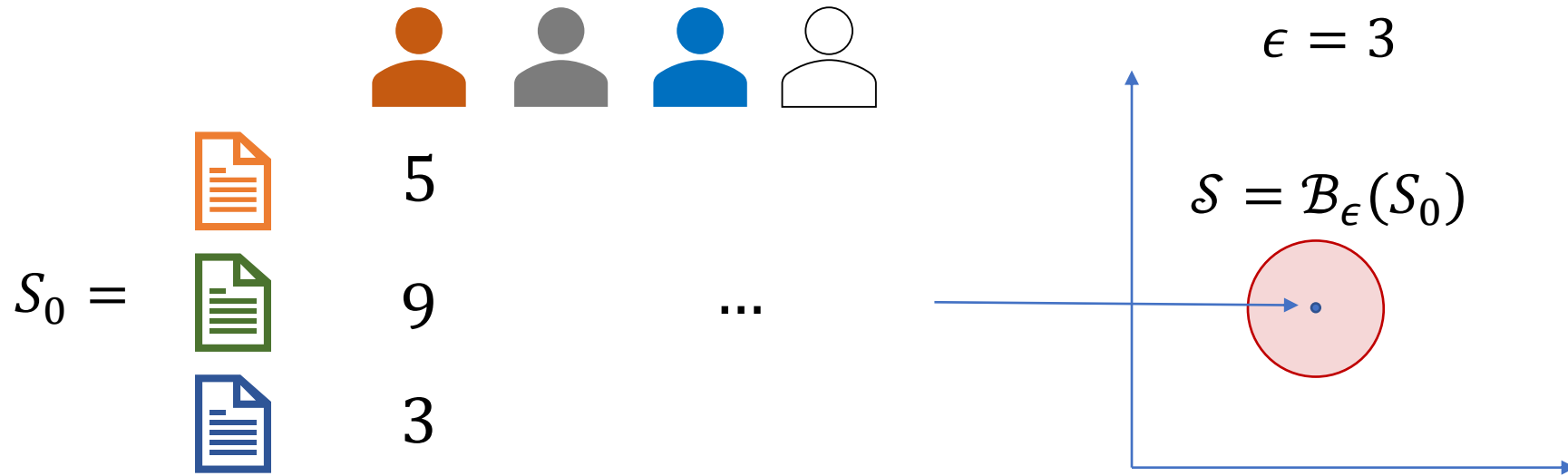
\mathcal{S} as Box Constraints



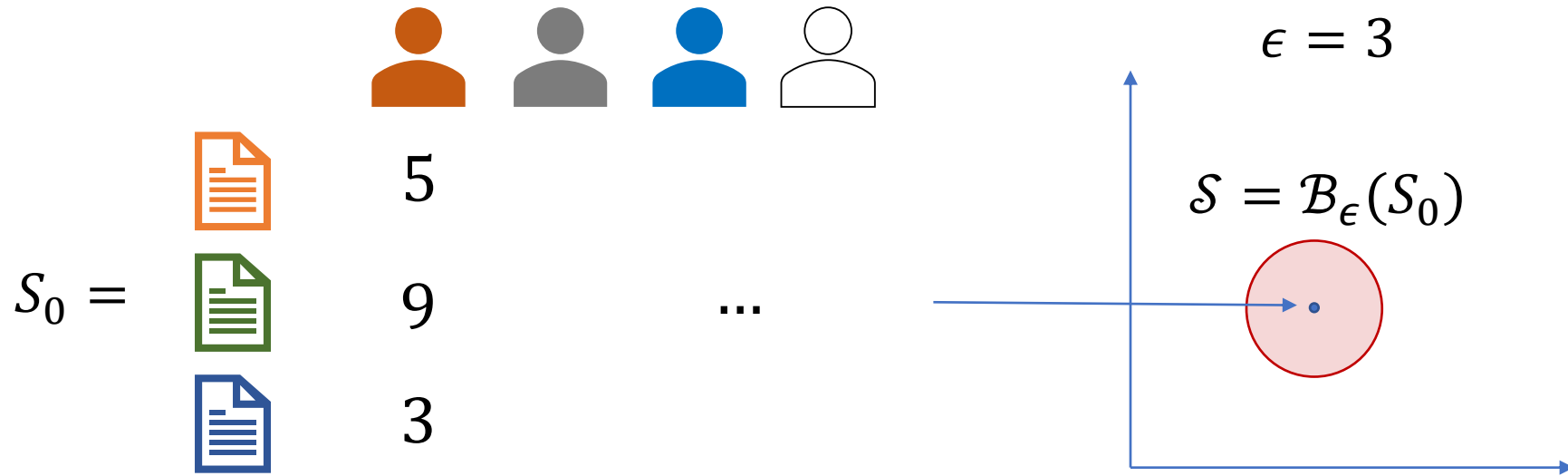
$$\underline{S} = \begin{array}{ccc} \begin{array}{c} \text{orange document icon} \\ \text{green document icon} \\ \text{blue document icon} \end{array} & \begin{array}{c} 4 \\ 7 \\ 1 \end{array} & \begin{array}{c} \\ \dots \\ \end{array} \end{array}$$

$$\begin{aligned} & \max_{A \in \mathcal{A}} \min_{S_A \in \mathcal{S}} USW(A, S_A) \\ &= \max_{A \in \mathcal{A}} USW(A, \underline{S}) \end{aligned}$$

\mathcal{S} as Spherical Noise



\mathcal{S} as Spherical Noise



Theorem:

$$\max_{A \in \mathcal{A}} \min_{S_A \in \mathcal{B}_\epsilon(S_0)} USW(A, S_A) = \max_{A \in \mathcal{A}} USW(A, S_0)$$

The welfare maximizer is robust to spherical noise.

\mathcal{S} as Spherical Noise

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$$\begin{aligned} \min_{S_A \in \mathcal{B}_\epsilon(S_0)} USW(A, S_A) &= \min_{x \in \mathcal{B}_\epsilon(0)} USW(A, S_0 + x) \\ &= USW(A, S_0) + \min_{x \in \mathcal{B}_\epsilon(0)} USW(A, x) \end{aligned}$$

\mathcal{S} as Spherical Noise

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Minimized when x points as far in the direction opposite A as possible:

$$x = -A \frac{\epsilon}{\|A\|_2}$$

\mathcal{S} as Spherical Noise

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Minimized when x points as far in the direction opposite A as possible:

$$x = -A \frac{\epsilon}{\|A\|_2}$$

Constant, due to constraints

AAAI's Affinity Scores

$$\text{Base Aggregated Score} = \begin{cases} \frac{1}{4}\text{TPMS} + \frac{1}{4}\text{ACL} + \frac{1}{2}\text{SAM} & \text{if all data is present} \\ \frac{1}{2}\text{ACL} + \frac{1}{2}\text{SAM} & \text{if TPMS is missing} \\ \frac{1}{2}\text{TPMS} + \frac{1}{2}\text{SAM} & \text{if ACL is missing} \\ \text{SAM} & \text{if both ACL and TPMS are missing} \end{cases}$$

Noise varies by score availability

Also modeled by elliptical noise

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Noise varies by score availability

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(Groupwise) Egalitarian Welfare

$$ESW(A, S) = \min_{i \in N} \sum_{j \in R} A_{ij} S_{ij}$$

If we have groups $G = \{G_1, G_2 \dots G_g\}$
(e.g. subject areas/tracks),
we define groupwise ESW:

$$ESW(A, S, G) = \min_{G_k \in G} \sum_{i \in G_k} \sum_{j \in R} A_{ij} S_{ij}$$

(Groupwise) Egalitarian Welfare

Fair to papers

$$ESW(A, S) = \min_{i \in N} \sum_{j \in R} A_{ij} S_{ij}$$

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Fair to subject areas

$$ESW(A, S, G) = \min_{G_k \in G} \sum_{i \in G_k} \sum_{j \in R} A_{ij} S_{ij}$$

Future Work

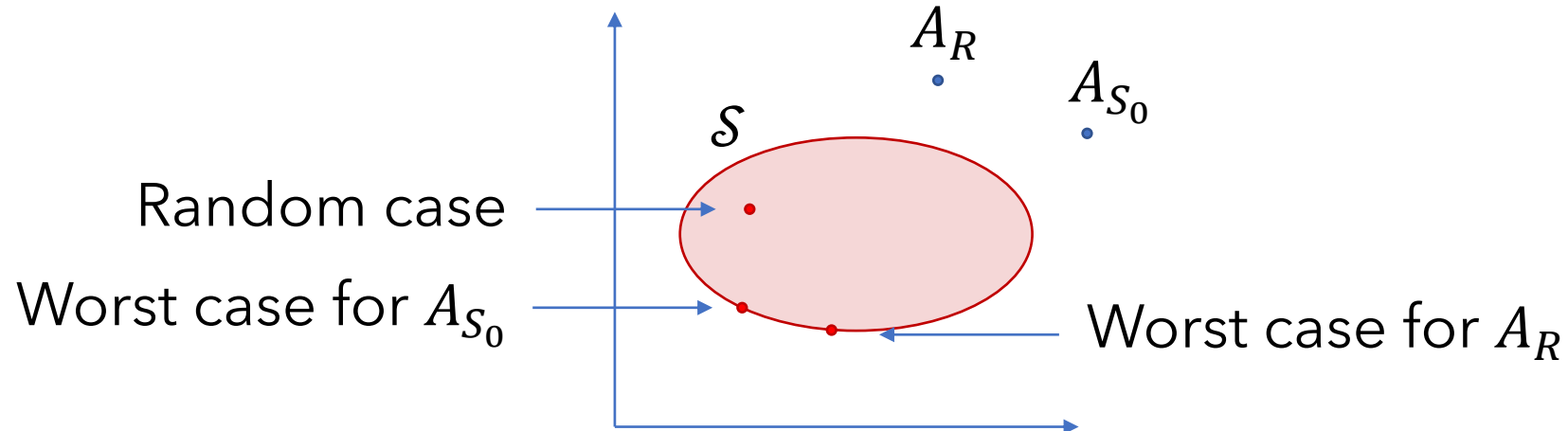
- Quantify integrality gaps
 - What is the impact of rounding?
- Closed-form solution for elliptical noise?
- More realistic experimental setup
 - Build document-similarity and bidding models with noise estimates
- Robust fairness
 - Envy term in the objective
 - Group-wise egalitarian welfare
- Other applications of robust allocation
 - Course allocation with inferred preferences, etc

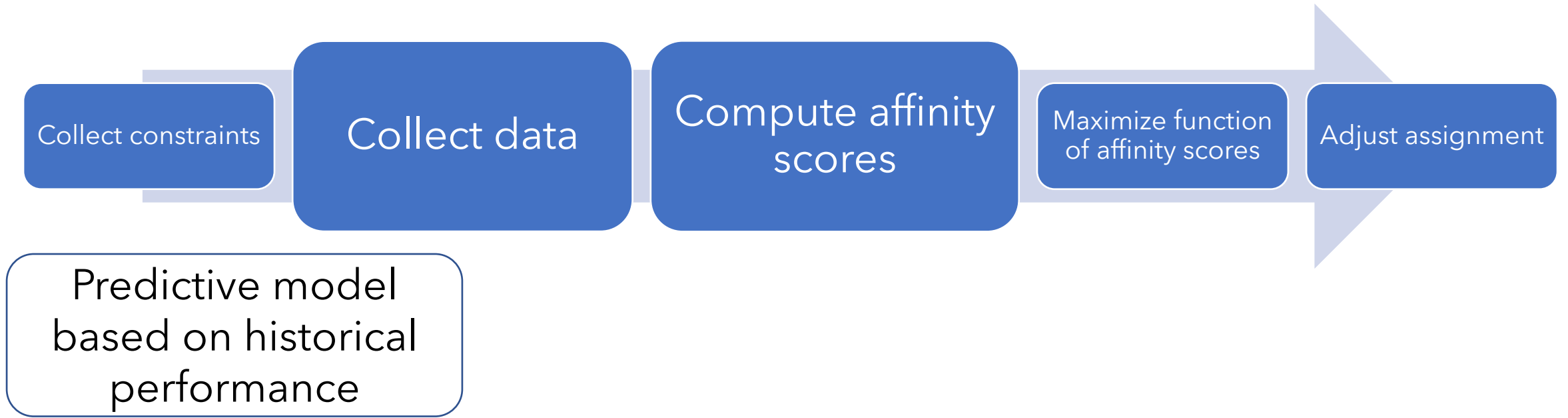
Measuring Robustness

Identify allocation A_{S_0} produced by solving for optimal under S_0

Concentrate noise around pairs assigned by A_{S_0} with parameter $1 - \alpha$

Compare worst-case and random-case welfare of A_{S_0} with robust solution A_R

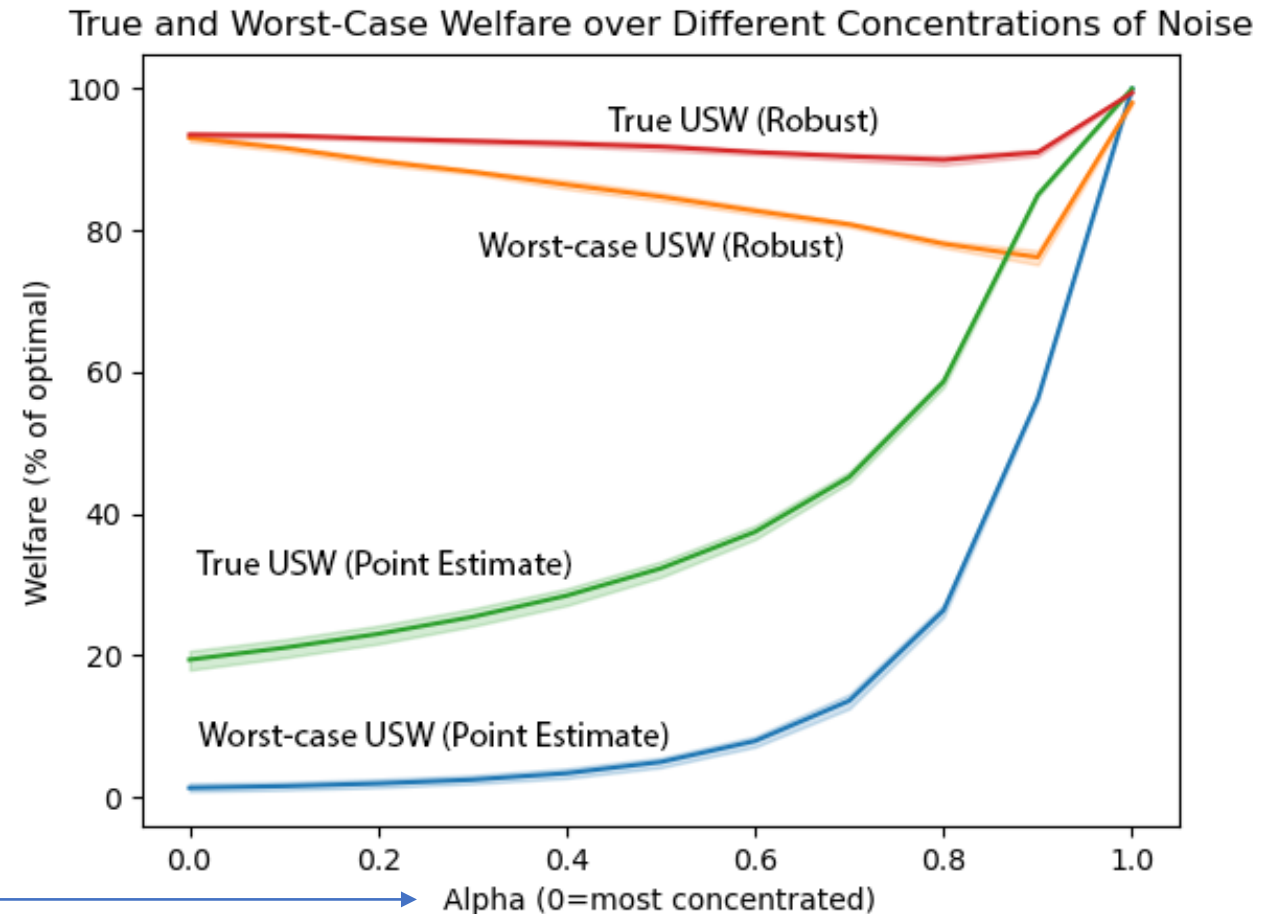




Measuring Robustness

Robust achieves higher worst-case USW and higher USW on randomly sampled true scores

α controls concentration of noise on pairs assigned by point estimate method



Computing the Subgradient

Initialize $A_1 = \max_{A \in \mathcal{A}} USW(A, S_0)$

For $t \in \{1, 2, \dots, T\}$:

$$S_t = \operatorname{argmin}_{S \in \mathcal{S}} USW(A_t, S)$$

$$A_{t+1} = A_t + \frac{1}{t+1} S_t$$

$$A_{t+1} = P_{\mathcal{A}}(A_{t+1})$$

Return rounded A_T

Projection onto
feasible set of
allocations \mathcal{A}

Round to an integer,
constraint-satisfying
allocation¹

1. Jecmen et al., NeurIPS 2020.